

1 Paraconsistent second order arithmetic $\mathbf{Z}_2^\#$ based on the paraconsistent logic $\overline{\text{LP}}_\omega^\#$ with infinite hierarchy levels of contradiction.Berry's and Richard's inconsistent numbers within $\mathbf{Z}_2^\#$.

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Abstract: In this paper paraconsistent second order arithmetic $\mathbf{Z}_2^\#$ with unrestricted comprehension scheme is proposed. We outline the development of certain portions of paraconsistent mathematics within paraconsistent second order arithmetic $\mathbf{Z}_2^\#$. In particular we defined infinite hierarchy Berry's and Richard's inconsistent numbers as elements of the paraconsistent field $\mathbb{R}^\#$.

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2 I.Introduction.

Let be \mathbf{Z}_2^* second order arithmetic [3]-[5] with second order language L_2 and with unrestricted comprehension scheme:

$$\exists X \forall n (n \in X \leftrightarrow \varphi(n, X))$$

It is known that second order arithmetic \mathbf{Z}_2^* is inconsistent from the well known standard construction named as Berry's and Richard's inconsistent numbers. Suppose that $F(n, X) \in L_2$ is a well-formed formula of second-order arithmetic \mathbf{Z}_2^* , i.e. formula which is arithmetical, which has one free set variable X and one free individual variable n . Suppose that $g(\exists X F(x, X)) \leq \mathbf{k}$, where $g(\exists X F(x, X))$ is a corresponding Gödel number. Let be $A_{\mathbf{k}}, \mathbf{k} \in \mathbb{N}$

the set of all positive consistent integers \bar{n} which can be defined under corresponding well-formed formula $F_{\bar{n}}(x, X)$, i.e. $\exists X_{\bar{n}} \forall m [F_{\bar{n}}(m, X_{\bar{n}}) \rightarrow m = \bar{n}]$, hence $\bar{n} \in A_{\mathbf{k}} \longleftrightarrow \exists X F_{\bar{n}}(\bar{n}, X_{\bar{n}})$.

$$\text{Thus } \forall n [n \in A_{\mathbf{k}} \longleftrightarrow \exists X_n F_n(n, X_n)], \quad (1.2)$$

$$g(\exists X_n F_n(x, X_n)) \leq \mathbf{k},$$

where $g(\exists X F(x, X))$ is a corresponding Gödel number. Since there are only finitely many of these \bar{n} , there must be a smallest positive integer $n_{\mathbf{k}} \in \mathbb{N} \setminus A_{\mathbf{k}}$ that does not belong to $A_{\mathbf{k}}$. But we just defined $n_{\mathbf{k}}$ in under corresponding well-

$$\text{formed formula } n_{\mathbf{k}} \in A_{\mathbf{k}} \longleftrightarrow \check{F}_{n_{\mathbf{k}}}(n_{\mathbf{k}}, A_{\mathbf{k}}), \quad (1.3)$$

$$\check{F}_{n_{\mathbf{k}}}(n_{\mathbf{k}}, A_{\mathbf{k}}) \longleftrightarrow n_{\mathbf{k}} = \min_{n \in \mathbb{N}} (\mathbb{N} \setminus A_{\mathbf{k}}).$$

Hence for a sufficiently Large \mathbf{k} such that: $g(\check{F}(n_{\mathbf{k}}, A_{\mathbf{k}})) \leq \mathbf{k}$ one obtain the contradiction: $(n_{\mathbf{k}} \in A_{\mathbf{k}}) \wedge (n_{\mathbf{k}} \notin A_{\mathbf{k}})$.

Within \mathbf{Z}_2^* , a consistent real number $x = \langle q_n : n \in \mathbb{N} \rangle \in \mathbb{R}$ is defined to be a Cauchy consistent sequence $\langle q_n | n \in \mathbb{N} \rangle$ of rational numbers, i.e., a consistent sequence of rational numbers $x = \langle q_n : n \in \mathbb{N}, q_n \in \mathbb{Q} \rangle$ such that

$$\forall \varepsilon (\varepsilon \in \mathbb{Q}) (\varepsilon > 0)$$

see **Definition 2.2.9.**

Let be $q_n \in \mathbb{Q}$ rational number with corresponding decimal representation $q_n = \{0, q_n(1) q_n(2) \dots q_n(n)\}$, $q_n(i) = 0, 1, 2, \dots, 9$, $i \leq n$, $x_k = \langle q_n^k \rangle = \langle q_n^k : n \in \mathbb{N}, q_n^k = \{0, q_n^k(1) q_n^k(2) \dots q_n^k(n)\} \rangle \in \mathbb{R}$ is a consistent real number which can be defined under corresponding well-formed formula (of second-order arithmetic \mathbf{Z}_2) $F_k(x)$, i.e. $\forall q (q \in \mathbb{Q}) [q \in \langle q_n^k \rangle \leftrightarrow \exists X F_k(q, X)]$. We denote real number x_k as k -th Richard's real number.

Let us consider Richard's real number $\mathfrak{R}_p = \langle \mathfrak{R}_n^p : n \in \mathbb{N} \rangle$ such that

$$\mathfrak{R}_n^p = 1 \leftrightarrow$$

$$\mathfrak{R}_n^p = 0 \leftrightarrow$$

Suppose that $\mathfrak{R}_p^p(p) \neq 1$, hence $\mathfrak{R}_p^p(p) = 1$. Thus $\mathfrak{R}_p^p(p) \neq q_p^p(p) \rightarrow \mathfrak{R}_p \neq x_p$. Suppose that $\mathfrak{R}_p^p(p) = 1$, hence $\mathfrak{R}_p^p(p) = 0$. Thus $\mathfrak{R}_p^p(p) \neq q_p^p(p) \rightarrow \mathfrak{R}_p \neq x_p$.

Hence for any Richard's real number x_k one obtain the contradiction $x_k \neq \mathfrak{R}_p^p(p)$. Thus the classical logical antinomy known as Richard-Berry paradox is combined with plausible assumptions formalizing certain sentences, to show that formalization of language leads to contradictions which trivialize the system \mathbf{Z}_2^* . In this paper paraconsistent second order arithmetic $\mathbf{Z}_2^\#$ with unrestricted comprehension scheme is proposed. We outline the development of certain portions of paraconsistent mathematics within paraconsistent second order arithmetic $\mathbf{Z}_2^\#$. In particular we defined infinite hierarchy Berry's and Richard's inconsistent numbers as elements of the paraconsistent field $\mathbb{R}^\#$.

3 II.Consistent second order arithmetic Z_2 .

4 II.1.System Z_2 .

In this section we briefly define Z_2 , the well known formal system of second order consistent arithmetic. For more detailed information concerning this system see [3]-[5].

The language L_2 of second order consistent arithmetic is a two-sorted language. This means that there are two distinct sorts of variables which are intended to range over two different kinds of object.

(1) **Variables:**

(1.1) Variables of the *first sort*: are known as consistent number variables, are denoted by i, j, k, m, n, \dots , and are intended to range over the set $\omega = \{0, 1, 2, \dots\}$ of all consistent natural numbers.

(1.2) Variables of the second sort are known as consistent set variables, are denoted by X, Y, Z, \dots , and are intended to range over all subsets of ω . The terms and formulas of the language of second order consistent arithmetic are as follows:

(2) Numerical terms are number variables, the constant symbols 0 and 1, and $t_1 + t_2$ and $t_1 \times t_2$ whenever t_1 and t_2 are numerical terms.

Here $(\cdot + \cdot)$ and $(\cdot \times \cdot)$ are binary operation symbols intended to denote addition and multiplication of natural numbers. (Numerical terms are intended to denote natural numbers.)

(3) Atomic formulas are:

(3.1) $t_1 = t_2$, $t_1 < t_2$,

(3.2) $t_1 \in X$ where t_1 and t_2 are numerical terms and X is any set variable. (The intended meanings of these respective atomic formulas are that t_1 equals t_2 , t_1 is less than t_2 , and t_1 is an element of X .)

(4) Formulas are built up from:

(4.1) atomic formulas by means of propositional connectives $\wedge, \vee, \neg, \rightarrow, \leftrightarrow$

(and, or, not, implies, if and only if),

(4.2) *consistent number* quantifiers $\forall n, \exists n$ (for all n , there exists n),

(4.3) *consistent set* quantifiers $\forall X, \exists X$ (for all X , there exists X).

(5) A sentence is a formula with no free variables.

Definition 2.1.1. (language of second order consistent arithmetic). L_2 is defined to be the language of second order consistent arithmetic as described above. In writing terms and formulas of L_2 , we shall use parentheses and brackets to indicate grouping, as is customary in mathematical logic textbooks. We shall also use some obvious abbreviations. For instance, $2 + 2 = 4$ stands for $(1 + 1) + (1 + 1) = ((1 + 1) + 1) + 1$, $(m + n)^2 \notin X$ stands for $\neg((m + n) \cdot (m + n) \in X)$, $s \leq t$ stands for $s < t \vee s = t$, and $\varphi \wedge \psi \wedge \theta$ stands for $(\varphi \wedge \psi) \wedge \theta$.

The semantics of the language L_2 are given by the following definition.

Definition 2.1.2. (L_2 -structures). A model for L_2 , also called a structure for L_2 or an L_2 -structure, is an ordered 7-tuple:

where $|\mathbf{M}|$ is a set which serves as the range of the number variables, \mathbf{S}_M is a set of subsets of $|\mathbf{M}|$ serving as the range of the set variables, $+_M$ \times_M are binary operations on $|\mathbf{M}|$, 0_M and 1_M are distinguished elements of $|\mathbf{M}|$, and $<_M$ is a binary relation on $|\mathbf{M}|$. We always assume that the sets $|\mathbf{M}|$ and \mathbf{S}_M are disjoint and nonempty. Formulas of L_2 are interpreted in \mathbf{M} in the obvious way.

In discussing a particular model \mathbf{M} as above, it is useful to consider formulas with parameters from $|\mathbf{M}| \cup \mathbf{S}_M$. We make the following slightly more general definition.

Definition 2.1.3. (parameters). Let \mathbf{B} be any subset of $|\mathbf{M}| \cup \mathbf{S}_M$. By a formula with parameters from \mathbf{B} we mean a formula of the extended language $L_2(\mathbf{B})$. Here $L_2(\mathbf{B})$ consists of L_2 augmented by new constant symbols corresponding to the elements of \mathbf{B} . By a sentence with parameters from \mathbf{B} we mean a sentence of $L_2(\mathbf{B})$, i.e., a formula of $L_2(\mathbf{B})$ which has no free variables.

In the language $L_2(|\mathbf{M}| \cup \mathbf{S}_M)$, constant symbols corresponding to elements of \mathbf{S}_M (respectively $|\mathbf{M}|$) are treated syntactically as unquantified set variables (respectively unquantified number variables). Sentences and formulas with parameters from $|\mathbf{M}| \cup \mathbf{S}_M$ are interpreted in \mathbf{M} in the obvious way.

Definition 2.1.4. A set $\mathbf{A} \subseteq |\mathbf{M}|$ is said to be definable over \mathbf{M} allowing parameters from \mathbf{B} if there exists a formula $\varphi(n)$ with parameters from \mathbf{B} and no free variables other than n such that $A = \{a \in |\mathbf{M}| : \mathbf{M} \models \varphi(a)\}$. Here $\mathbf{M} \models \varphi(a)$ means that \mathbf{M} satisfies $\varphi(a)$, i.e., $\varphi(a)$ is true in \mathbf{M} .

Definition 2.1.5. The *intended model* for L_2 is of course the model $(\omega, \mathbf{P}(\omega), +, \times, 0, 1, <)$ where ω is the set of natural numbers, $\mathbf{P}(\omega)$ is the set of all subsets of ω , and $+, \cdot, 0, 1, <$ are as usual.

By an ω -model we mean an L_2 -structure of the form $(\omega, \mathbf{S}, +, \cdot, 0, 1, <)$ where $\emptyset \neq \mathbf{S} \subseteq \mathbf{P}(\omega)$. Thus an ω -model differs from the intended model only by having a possibly smaller collection \mathbf{S} of sets to serve as the range of the set variables. We sometimes speak of the ω -model \mathbf{S} when we really mean the ω -model $(\omega, \mathbf{S}, +, \cdot, 0, 1, <)$.

Definition 2.1.6. (second order arithmetic \mathbf{Z}_2). The axioms of second order

arithmetic \mathbf{Z}_2 consist of the universal closures of the following L_2 -formulas:

(i) basic axioms:

(i.1) $n + 1 \neq 0$,

(i.2) $m + 1 = n + 1 \rightarrow m = n$,

(i.3) $m + 0 = m$,

(i.4) $m + (n + 1) = (m + n) + 1$,

(i.5) $m \times 0 = 0$,

(i.6) $m \times (n + 1) = (m \times n) + m$,

(i.7) $\neg(m < 0)$,

(i.8) $m < n + 1 \leftrightarrow (m < n \vee m = n)$.

(ii) induction axiom:

$$\mathbf{M} = (|\mathbf{M}|, \mathbf{S}_M, (\cdot +_M \cdot),$$

and

$$(0 \in X \wedge \forall n(n \in X \rightarrow n + 1 \in X)) \rightarrow \forall n(n \in X)$$

(iii) comprehension scheme:

$$\exists X \forall n(n \in X \leftrightarrow \varphi(n))$$

where $\varphi(n)$ is any formula of L_2 in which X does not occur freely.

5 II.2.Consistent Mathematics Within \mathbf{Z}_2 .

We now outline the development of certain portions of ordinary mathematics within \mathbf{Z}_2 .

Definition 2.2.1. If X and Y are set variables, we use $X = Y$ and $X \subseteq Y$ as

abbreviations for the formulas $\forall n(n \in X \leftrightarrow n \in Y)$ and $\forall n(n \in X \rightarrow n \in Y)$

respectively.

Definition 2.2.2. Within \mathbf{Z}_2 , we define \mathbb{N} to be the unique set X such that $\forall n(n \in X)$.

Definition 2.2.3. For $X, Y \subseteq \mathbb{N}$, a consistent function $f : X \rightarrow Y$ is defined to be a

consistent set $f \subseteq X \times Y$ such that for all $m \in X$ there is exactly one $n \in Y$ such

that $(m, n) \in f$. For $m \in X$, $f(m)$ is defined to be the unique n such that $(m, n) \in f$.

The usual properties of such functions can be proved in \mathbf{Z}_2 .

Definition 2.2.4. (consistent primitive recursion). This means that, given

$f : X \rightarrow Y$ and $g : \mathbb{N} \times X \times Y \rightarrow Y$, there is a unique $h : \mathbb{N} \times X \rightarrow Y$ defined by

$$h(0, m) = f(m),$$

$$h(n + 1, m) = g(n, m, h(n, m)) \text{ for all } n \in \mathbb{N} \text{ and } m \in X.$$

The existence of h is proved by arithmetical comprehension, and the uniqueness of h is proved by arithmetical induction.

In particular, we have the exponential function $\exp(m, n) = m^n$, defined by $m^0 = 1, m^{n+1} = m^n \times m$ for all $m, n \in \mathbb{N}$. The usual properties of the exponential

function can be proved in \mathbf{Z}_2 .

The consistent natural number system is essentially already given to us by the

language L_2 and axioms of \mathbf{Z}_2 . Thus, within \mathbf{Z}_2 , a consistent natural number is

defined to be an element of \mathbb{N} , and the natural number system is defined to be

the structure $\mathbb{N}, +_{\mathbb{N}}, \times_{\mathbb{N}}, 0_{\mathbb{N}}, 1_{\mathbb{N}}, <_{\mathbb{N}}, =_{\mathbb{N}}$, where $+_{\mathbb{N}} : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ is defined by $m +_{\mathbb{N}} n = m + n$, etc. Thus for instance $+_{\mathbb{N}}$ is the set of triples $((m, n), k) \in (\mathbb{N} \times \mathbb{N}) \times \mathbb{N}$ such that $m + n = k$. The existence of this set follows

from the arithmetical comprehension.

In a standard manner, we can define within \mathbf{Z}_2 the set \mathbb{Z} of consistent integers

and the set of consistent rational numbers: \mathbb{Q} .

Definition 2.2.5. (consistent rational numbers \mathbb{Q}) Let $\mathbb{Z}^+ = \{a \in \mathbb{Z} : 0 <_{\mathbb{Z}} a\}$ be

the set of positive consistent integers, and let $\equiv_{\mathbb{Q}}$ be the equivalence relation

on $\mathbb{Z} \times \mathbb{Z}^+$ defined by $(a, b) \equiv_{\mathbb{Q}} (c, d)$ if and only if $a \times_{\mathbb{Z}} d = b \times_{\mathbb{Z}} c$. Then \mathbb{Q} is

defined to be the set of all $(a, b) \in \mathbb{Z} \times \mathbb{Z}^+$ such that (a, b) is the $<_{\mathbb{N}}$ -minimum

element of its $\equiv_{\mathbb{Q}}$ -equivalence class. Operations $+_{\mathbb{Q}}, -_{\mathbb{Q}}, \times_{\mathbb{Q}}$ on \mathbb{Q} are defined

by:

$(a, b) +_{\mathbb{Q}} (c, d) \equiv_{\mathbb{Q}} (a \times_{\mathbb{Z}} d +_{\mathbb{Z}} b \times_{\mathbb{Z}} c, b \times_{\mathbb{Z}} d)$, $-_{\mathbb{Q}}(a, b) \equiv_{\mathbb{Q}} (-_{\mathbb{Z}} a, b)$, and

$(a, b) \times_{\mathbb{Q}} (c, d) \equiv_{\mathbb{Q}} (a \times_{\mathbb{Z}} c, b \times_{\mathbb{Z}} d)$. We let $0_{\mathbb{Q}} \equiv_{\mathbb{Q}} (0_{\mathbb{Z}}, 1_{\mathbb{Z}})$ and $1_{\mathbb{Q}} \equiv_{\mathbb{Q}} (1_{\mathbb{Z}}, 1_{\mathbb{Z}})$,

and we define a binary relation $<_{\mathbb{Q}}$ on \mathbb{Q} by letting $(a, b) <_{\mathbb{Q}} (c, d)$ if and only

if $a \times_{\mathbb{Z}} d <_{\mathbb{Z}} b \times_{\mathbb{Z}} c$. Finally $=_{\mathbb{Q}}$ is the identity relation on \mathbb{Q} . We can then prove

within \mathbf{Z}_2 that the rational number system $\mathbb{Q}, +_{\mathbb{Q}}, -_{\mathbb{Q}}, \times_{\mathbb{Q}}, 0_{\mathbb{Q}}, 1_{\mathbb{Q}}, <_{\mathbb{Q}}, =_{\mathbb{Q}}$ has the

usual properties of an ordered field, etc.

We make the usual identifications whereby \mathbb{N} is regarded as a subset of \mathbb{Z} and

\mathbb{Z} is regarded as a subset of \mathbb{Q} . Namely $m \blacksquare \mathbb{N}$ is identified with $(m, 0) \in \mathbb{Z}$,

and $a \in \mathbb{Z}$ is identified with $(a, 1_{\mathbb{Z}}) \in \mathbb{Q}$. We use $+$ ambiguously to denote $+_{\mathbb{N}}, +_{\mathbb{Z}}$,

or $+_{\mathbb{Q}}$ and similarly for $-$, \times , 0 , 1 , $<$. For $q, r \in \mathbb{Q}$ we write $q - r = q + (-r)$, and if

$r \neq 0$, $q/r =$ the unique $q' \in \mathbb{Q}$ such that $q = q' \times r$. The function $\exp(q, a) = q^a$ for

$q \in \mathbb{Q} \setminus \{0\}$ and $a \in \mathbb{Z}$ is obtained by primitive recursion in the obvious way.

Definition 2.2.6. The absolute value function $|\cdot| : \mathbb{Q} \rightarrow \mathbb{Q}$ is defined by

$$|q| = q$$

if $q \geq 0$, $-q$ otherwise.

Definition 2.2.7. An consistent sequence of rational numbers is defined to be a

consistent function $f : \mathbb{N} \rightarrow \mathbb{Q}$.

We denote such a sequence as $\langle q_n : n \in \mathbb{N} \rangle$, or simply $\langle q \rangle_n$, where $q_n = f(n)$.

Definition 2.2.8. A double consistent sequence of rational numbers

is a consistent function $f : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{Q}$, denoted $\langle q_{mn} : m, n \in \mathbb{N} \rangle$ or simply $\langle q_{mn} \rangle$,

where $q_{mn} = f(m, n)$.

Definition 2.2.9. (consistent real numbers). Within \mathbf{Z}_2 , a consistent real number

is defined to be a Cauchy consistent sequence of rational numbers, i.e., a

consistent sequence of rational numbers $x = \langle q_n : n \in \mathbb{N} \rangle$ such that

Definition 2.2.10. If $x = q_n$ and $y = q'_n$ are consistent real numbers, we

write $x =_{\mathbb{R}} y$ to mean that $\lim_n |q_n - q'_n| = 0$, i.e.,

$$\forall \varepsilon (\varepsilon \in \mathbb{Q}) (\varepsilon > 0 \rightarrow \exists m \forall n (m < n \rightarrow |q_n - q'_n| < \varepsilon)).$$

and we write $x <_{\mathbb{R}} y$ to mean that

$$\exists \varepsilon (\varepsilon > 0 \wedge \exists m \forall n (m < n \rightarrow q_n + \varepsilon < q'_n)).$$

Also $x +_{\mathbb{R}} y = \langle q_n + q'_n \rangle$, $x \times_{\mathbb{R}} y = \langle q_n \times q'_n \rangle$, $-_{\mathbb{R}} x = \langle -q_n \rangle$, $0_{\mathbb{R}} = \langle 0 \rangle$, $1_{\mathbb{R}} = \langle 1 \rangle$.

We use \mathbb{R} to denote the set of all *consistent real numbers*. Thus $x \in \mathbb{R}$ means that x is a *consistent real number*. (Formally, we cannot speak of the

set \mathbb{R} within the language of second order arithmetic, since it is a set of sets.)

We shall usually omit the subscript \mathbb{R} in $+_{\mathbb{R}}, -_{\mathbb{R}}, \times_{\mathbb{R}}, 0_{\mathbb{R}}, 1_{\mathbb{R}}, <_{\mathbb{R}}, =_{\mathbb{R}}$.

Thus the *consistent real number system* consists of $\mathbb{R}, +, -, \times, 0, 1, <, =$. We shall

sometimes identify a consistent rational number $q \in \mathbb{Q}$ with the corresponding

consistent real number $x_q = \langle q \rangle$.

Within \mathbf{Z}_2 one can prove that the real number system has the usual properties of an *consistent Archimedean ordered field*, etc. The complex

consistent numbers can be introduced as usual as pairs of real numbers.

Within \mathbf{Z}_2 , it is straightforward to carry out the proofs of all the basic results in

real and complex linear and polynomial algebra. For example, the fundamental

theorem of algebra can be proved in \mathbf{Z}_2 .

Definition 2.2.11. A consistent sequence of real numbers is defined to be a

double consistent sequence of rational numbers $\langle q_{mn} : m, n \in \mathbb{N} \rangle$ such that for each m , $\langle q_{mn} : n \in \mathbb{N} \rangle$ is a consistent real number. Such a sequence of real

numbers is denoted $\langle x_m : m \in \mathbb{N} \rangle$, where $x_m = \langle q_{mn} : n \in \mathbb{N} \rangle$. Within \mathbf{Z}_2 we can

prove that every bounded consistent sequence of real numbers has a *consistent least upper bound*. This is a very useful completeness property of

the consistent real number system. For instance, it implies that an infinite series

of positive terms is convergent if and only if the finite partial sums are bounded.

We now turn of certain portions of consistent abstract algebra within \mathbf{Z}_2 .

Because of the restriction to the language L_2 of second order arithmetic, we

cannot expect to obtain a good general theory of arbitrary (countable and uncountable) algebraic structures. However, we can develop countable algebra, i.e., the theory of countable algebraic structures, within \mathbf{Z}_2 .

Definition 2.2.12. A countable consistent commutative ring is defined within \mathbf{Z}_2

to be a consistent structure $\mathbf{R}, +_{\mathbf{R}}, -_{\mathbf{R}}, \times_{\mathbf{R}}, 0_{\mathbf{R}}, 1_{\mathbf{R}}$, where $\mathbf{R} \subseteq \mathbb{N}$, $+_{\mathbf{R}} : \mathbf{R} \times \mathbf{R} \rightarrow \mathbf{R}$, etc., and the usual commutative ring axioms are assumed.

(We include $0 \neq 1$ among those axioms.) The subscript \mathbf{R} is usually omitted.

An ideal in \mathbf{R} is a set $I \subseteq \mathbf{R}$ such that $a \in I$ and $b \in I$ imply $a + b \in I$; $a \in I$ and

$r \in \mathbf{R}$ imply $a \times r \in I$, and $0 \in I$ and $1 \notin I$. We define an equivalence relation

$=_I$ on \mathbf{R} by $r =_I s$ if and only if $r - s \in I$. We let \mathbf{R}/I be the set of $r \in \mathbf{R}$ such

that r is the $<_{\mathbb{N}}$ -minimum element of its equivalence class under $=_I$. Thus \mathbf{R}/I

consists of one element of each $=_I$ -equivalence class of elements of \mathbf{R} . With

the appropriate operations, \mathbf{R}/I becomes a countable commutative ring, the

quotient ring of \mathbf{R} by I .

The ideal I is said to be prime if \mathbf{R}/I is an integral domain, and maximal if \mathbf{R}/I

is a field.

Next we indicate how some basic concepts and results of analysis and topology can be developed within \mathbf{Z}_2 .

Definition 2.2.13. Within \mathbf{Z}_2 , a complete separable consistent metric space is a

nonempty set $A \subseteq \mathbb{N}$ together with a function $d : A \times A \rightarrow \mathbb{R}$ satisfying $d(a, a) = 0$, $d(a, b) = d(b, a) \geq 0$, and $d(a, c) \leq d(a, b) + d(b, c)$ for all $a, b, c \in A$.

(Formally, d is a consistent sequence of real numbers, indexed by $A \times A$.) We

define a point \widehat{A} of the complete separable metric space \widehat{A} to be a sequence $x = \langle a_n : n \in \mathbb{N} \rangle$, $a_n \in A$, satisfying $\forall \varepsilon (\varepsilon \in \mathbb{R}) (\varepsilon > 0 \rightarrow \exists m \forall n (m < n \rightarrow d(a_m, a_n) < \varepsilon))$.

The pseudometric d is extended from A to \widehat{A} by

$$d(x, y) = \lim_{n \rightarrow \infty} d(a_n, b_n) \quad (2.2.5)$$

where $x = \langle a_n : n \in \mathbb{N} \rangle$ and $y = \langle b_n : n \in \mathbb{N} \rangle$. We write $x = y$ if and only if $d(x, y) = 0$. For example, $\mathbb{R} = \widehat{\mathbb{Q}}$ under the metric $d(q, q') = |q - q'|$.

Definition 2.2.14. (consistent continuous functions). Within \mathbf{Z}_2 , if \widehat{A} and \widehat{B} are

complete separable metric spaces, a consistent continuous function $\phi : \widehat{A} \rightarrow \widehat{B}$

is a set $\Phi \subseteq A \times \mathbb{Q}^+ \times B \times \mathbb{Q}^+$ satisfying the following coherence conditions:

1. $[(a, r, b, s) \in \Phi] \wedge [(a, r, b', s') \in \Phi] \dashrightarrow d(b, b') < s + s'$;
2. $[(a, r, b, s) \in \Phi] \wedge [d(b, b') + s < s'] \dashrightarrow (a, r, b', s') \in \Phi$ (2.2.6)
3. $[(a, r, b, s) \in \Phi] \wedge [d(a, a') + r' < r] \dashrightarrow (a', r', b, s) \in \Phi$

6 III.Paraconsistent second order arithmetic $\mathbf{Z}_2^\#$.

7 III.1.Paraconsistent system $\mathbf{Z}_2^\#$.

In this section we define \mathbf{Z}_2 , the formal system of second order paraconsistent arithmetic based on the paraconsistent logic $\overline{\mathbf{LP}}_\omega^\#$ [1] with infinite hierarchy levels of contradiction. For detailed information concerning paraconsistent logic $\mathbf{LP}_\omega^\#$ see [2].

Definition 3.1.1. For arbitrary binary inconsistent relation $(\cdot \sim \cdot)$ we define:

$$\begin{aligned} a \sim_{w,(0)} b &\triangleq (a \sim_w b)^{(0)}, \dots, a \sim_{w,(n)} b \triangleq (a \sim_w b)^{(n)}; \\ a \sim_{w,[0]} b &\triangleq (a \sim_w b)^{[0]}, \dots, a \sim_{w,[n]} b \triangleq (a \sim_w b)^{[n]}; n \in \mathbb{N}. \end{aligned} \quad (3.1.1)$$

$$\mathbb{N} = \{0, 1, 2, \dots\}$$

Definition 3.1.2. In particular we define:

$$\begin{aligned}
a =_{w,(0)} b &\triangleq (a =_w b)^{(0)}, \dots, a =_{w,(n)} b \triangleq (a =_w b)^{(n)}; \\
a =_{w,[0]} b &\triangleq (a =_w b)^{[0]}, \dots, a =_{w,[n]} b \triangleq (a =_w b)^{[n]}; n \in \mathbb{N}. \quad (3.1.2) \\
a <_{w,(0)} b &\triangleq (a <_w b)^{(0)}, \dots, a <_{w,(n)} b \triangleq (a <_w b)^{(n)}; \\
a <_{w,[0]} b &\triangleq (a <_w b)^{[0]}, \dots, a <_{w,[n]} b \triangleq (a <_w b)^{[n]}; n \in \mathbb{N}
\end{aligned}$$

The language $L_2^\#$ of second order paraconsistent arithmetic $\mathbf{Z}_2^\#$ is a two-sorted paraconsistent language. This means that there are two distinct sorts of variables which are intended to range over two different kinds of object.

(1) Variables:

(1.1) Variables of the *first sort*: are denoting *consistent and inconsistent* number

variables, are denoted as in classical case by i, j, k, m, n, \dots , and are intended to

range over the set $\mathbb{N}^\# \supsetneq \mathbb{N} = \{0, 1, 2, \dots\}$ of all *consistent and inconsistent*

natural numbers.

(1.2) Variables of the *second sort*: are denoting *consistent and inconsistent* set

variables, are denoted by X, Y, Z, \dots , and are intended to range over all sub-sets

of $\mathbb{N}^\#$. The terms and formulas of the language $L_2^\#$ of second order paraconsistent arithmetic $\mathbf{Z}_2^\#$ are as follows:

(2) Numerical terms are number variables, the constant symbols

$0_s, 0_w, 0_{w,(i)}, 0_{w,[i]}, 1_s, 1_w, 1_{w,(i)}, 1_{w,[i]}, i \in \mathbb{N}$ and $\mathbf{t}_1 + \mathbf{t}_2$ $\mathbf{t}_1 \times \mathbf{t}_2$ whenever \mathbf{t}_1 and \mathbf{t}_2 are

numerical terms in general.

Here $(\cdot + \cdot)$ and $(\cdot \times \cdot)$ are binary operation symbols intended to denote

addition and multiplication of *consistent and inconsistent* natural numbers.

(Numerical terms are intended to denote *consistent and inconsistent* natural numbers.)

(3) Atomic formulas are:

(3.1) $\mathbf{t}_1 =_s \mathbf{t}_2, \mathbf{t}_1 <_s \mathbf{t}_2, \mathbf{t}_1 =_w \mathbf{t}_2, \mathbf{t}_1 <_w \mathbf{t}_2,$

(3.2) $\mathbf{t}_1 =_{w,(0)} \mathbf{t}_2, \mathbf{t}_1 =_{w,(1)} \mathbf{t}_2, \dots, \mathbf{t}_1 =_{w,(n)} \mathbf{t}_2, n \in \mathbb{N},$

(3.3) $\mathbf{t}_1 =_{w,[0]} \mathbf{t}_2, \mathbf{t}_1 =_{w,[1]} \mathbf{t}_2, \dots, \mathbf{t}_1 =_{w,[n]} \mathbf{t}_2, n \in \mathbb{N},$

(3.4) $\mathbf{t}_1 <_{w,(0)} \mathbf{t}_2, \mathbf{t}_1 <_{w,(1)} \mathbf{t}_2, \dots, \mathbf{t}_1 <_{w,(n)} \mathbf{t}_2, n \in \mathbb{N},$

(3.5) $\mathbf{t}_1 <_{w,[0]} \mathbf{t}_2, \mathbf{t}_1 <_{w,[1]} \mathbf{t}_2, \dots, \mathbf{t}_1 <_{w,[n]} \mathbf{t}_2, n \in \mathbb{N},$

(3.6) $\mathbf{t}_1 \in_s X, \mathbf{t}_1 \in_w X,$

(3.7) $\mathbf{t}_1 \in_{w,(0)} \mathbf{t}_2, \mathbf{t}_1 \in_{w,(1)} \mathbf{t}_2, \dots, \mathbf{t}_1 \in_{w,(n)} \mathbf{t}_2, n \in \mathbb{N},$

(3.8) $\mathbf{t}_1 \in_{w,[0]} \mathbf{t}_2, \mathbf{t}_1 \in_{w,[1]} \mathbf{t}_2, \dots, \mathbf{t}_1 \in_{w,[n]} \mathbf{t}_2, n \in \mathbb{N},$

where \mathbf{t}_1 and \mathbf{t}_2 are numerical terms and X is any set variable.

The intended meanings of these respective atomic formulas are that:

- (3.1) t_1 equals t_2 in a *strong consistent sense*,
 t_1 is less than t_2 , in a *strong consistent sense*,
 t_1 equals t_2 in a *weak inconsistent sense*,
 t_1 is less than t_2 , in a *weak inconsistent sense*;
- (3.2) t_1 equals t_2 in a *weak inconsistent sense with rank = 0*,
 t_1 equals t_2 in a *weak inconsistent sense with rank = 1, 2, ...*,
 t_1 equals t_2 in a *weak consistent sense with rank = n*, $n \in \mathbb{N}$;
- (3.3) t_1 equals t_2 in a *strictly inconsistent sense with rank = 0*,
 t_1 equals t_2 in a *strictly inconsistent sense with rank = 1, ...*,
 t_1 equals t_2 in a *strictly inconsistent sense with rank = n*, $n \in \mathbb{N}$;
- (3.4) t_1 is less than t_2 , in a *weak inconsistent sense with rank = 0*,
 t_1 is less than t_2 , in a *weak inconsistent sense with rank = 1, 2, ...*,
 t_1 is less than t_2 , in a *weak consistent sense with rank = n*, $n \in \mathbb{N}$;
- (3.5) t_1 is less than t_2 , in a *strictly inconsistent sense with rank = 0*,
 t_1 is less than t_2 , in a *strictly inconsistent sense with rank = 1, 2, ...*,
 t_1 is less than t_2 , in a *strictly inconsistent sense with rank = n*, $n \in \mathbb{N}$;
- (3.6) t_1 is an element of X in a *strong consistent sense*,
 t_1 is an element of X in a *weak inconsistent sense*,
- (3.7) t_1 is an element of X in a *weak inconsistent sense with rank = 0*,
 t_1 is an element of X in a *weak inconsistent sense with rank = 1, 2, ...*,
 t_1 is an element of X in a *weak inconsistent sense with rank = n*;
- (3.8) t_1 is an element of X in a *strictly inconsistent sense with rank = 0*,
 t_1 is an element of X in a *strictly inconsistent sense with rank = 1, 2, ...*,
 t_1 is an element of X in a *strictly inconsistent sense with rank = n*, $n \in \mathbb{N}$;
- (4) Formulas are built up from:
 - (4.1) atomic formulas by means of propositional connectives $\wedge, \vee, \neg, \rightarrow, \leftrightarrow$
 - (and, or, not, implies, if and only if),
 - (4.2) *consistent and inconsistent number quantifiers* $\forall n, \exists n$ (for all n , there exists n),
 - (4.3) *consistent and inconsistent set quantifiers* $\forall X, \exists X$ (for all X , there exists X),
 - (4.4) operators $(\cdot)^{(0)}, (\cdot)^{(1)}, \dots, (\cdot)^{(n)}, (\cdot)^{[0]}, (\cdot)^{[1]}, \dots, (\cdot)^{[n]}$, $n \in \mathbb{N}$.
- (5) A sentence is a formula with no free variables.

Definition 3.1.3. (language of second order inconsistent arithmetic). $L_2^\#$ is defined to be the language of second order inconsistent arithmetic as described above. In writing terms and formulas of $L_2^\#$, we shall use parentheses and brackets to indicate grouping, as is customary in mathematical logic textbooks.

The semantics of the paraconsistent language $L_2^\#$ are given by the following definition.

Definition 3.1.4. (paraconsistent $L_2^\#$ -structures). A paraconsistent model for $L_2^\#$, also called a paraconsistent structure for $L_2^\#$ or an paraconsistent $L_2^\#$ -

$$\mathbf{M}_{\text{inc}} = \check{\mathbf{M}} = \left\{ |\check{\mathbf{M}}|, \mathbf{S}_{\check{\mathbf{M}}}, (\cdot +_{\check{\mathbf{M}}}\cdot), (\cdot \times_{\check{\mathbf{M}}}\cdot), 0_{\mathbf{s}}^{\check{\mathbf{M}}}, 1_{\mathbf{s}}^{\check{\mathbf{M}}}, 0_w^{\check{\mathbf{M}}}, 1_w^{\check{\mathbf{M}}}, \{0_{w,[n]}^{\check{\mathbf{M}}}\}_{n \in \mathbb{N}}, \{1_{w,(n)}^{\check{\mathbf{M}}}\}_{n \in \mathbb{N}}, \{1_{w,[n]}^{\check{\mathbf{M}}}\}_{n \in \mathbb{N}}, (\cdot =_{\mathbf{s}}^{\check{\mathbf{M}}}\cdot), (\cdot =_{w,[n]}^{\check{\mathbf{M}}}\cdot), (\cdot <_{\mathbf{s}}^{\check{\mathbf{M}}}\cdot), (\cdot <_{w,(n)}^{\check{\mathbf{M}}}\cdot), (\cdot <_{w,[n]}^{\check{\mathbf{M}}}\cdot) \right\}$$

structure, is an ordered 19-tuple:

$$\left\{ (\cdot =_{w,(n)}^{\check{\mathbf{M}}}\cdot) \right\}_{n \in \mathbb{N}}, \left\{ (\cdot =_{w,[n]}^{\check{\mathbf{M}}}\cdot) \right\}_{n \in \mathbb{N}},$$

$$(\cdot <_{\mathbf{s}}^{\check{\mathbf{M}}}\cdot), (\cdot <_w^{\check{\mathbf{M}}}\cdot), \left\{ (\cdot <_{w,(n)}^{\check{\mathbf{M}}}\cdot) \right\}_{n \in \mathbb{N}}, \left\{ (\cdot <_{w,[n]}^{\check{\mathbf{M}}}\cdot) \right\}_{n \in \mathbb{N}}$$

$$\mathbb{N} = \{0, 1, 2, \dots\},$$

where $|\mathbf{M}|_{\text{inc}} = |\check{\mathbf{M}}|$ is an *inconsistent set* which serves as the range of the consistent and inconsistent number variables, $\mathbf{S}_{\check{\mathbf{M}}}$ is a set of subsets of $|\check{\mathbf{M}}|$ serving as the range of the set variables, $+_{\check{\mathbf{M}}}$ and $\times_{\check{\mathbf{M}}}$ are binary operations on $|\check{\mathbf{M}}|$, $0_{\check{\mathbf{M}}} \triangleq \{0_{\mathbf{s}}^{\check{\mathbf{M}}}, 0_w^{\check{\mathbf{M}}}, 0_{w,(n)}^{\check{\mathbf{M}}}, 0_{w,[n]}^{\check{\mathbf{M}}}\}$ and $1_{\check{\mathbf{M}}} \triangleq \{1_{\mathbf{s}}^{\check{\mathbf{M}}}, 1_w^{\check{\mathbf{M}}}, 1_{w,(n)}^{\check{\mathbf{M}}}, 1_{w,[n]}^{\check{\mathbf{M}}}\}$, $n \in \mathbb{N}$ are distinguished elements of $|\check{\mathbf{M}}|$, and $(\cdot =_{\mathbf{s}}^{\check{\mathbf{M}}}\cdot), (\cdot <_{\mathbf{s}}^{\check{\mathbf{M}}}\cdot)$ is a binary *strongly consistent relations* on $|\check{\mathbf{M}}|$, $(\cdot =_w^{\check{\mathbf{M}}}\cdot), (\cdot =_{w,(n)}^{\check{\mathbf{M}}}\cdot), (\cdot <_w^{\check{\mathbf{M}}}\cdot), (\cdot <_{w,(n)}^{\check{\mathbf{M}}}\cdot), (\cdot <_{w,[n]}^{\check{\mathbf{M}}}\cdot)$, $n \in \mathbb{N}$ is a binary *weakly inconsistent relations* on $|\check{\mathbf{M}}|$, $(\cdot =_{w,[n]}^{\check{\mathbf{M}}}\cdot), (\cdot <_{w,[n]}^{\check{\mathbf{M}}}\cdot)$, $n \in \mathbb{N}$ is a binary *inconsistent relations* on $|\check{\mathbf{M}}|$. We always assume that the *inconsistent sets* $|\check{\mathbf{M}}|$ and $\mathbf{S}_{\check{\mathbf{M}}}$ are disjoint and nonempty. Formulas of $L_2^\#$ are interpreted in *inconsistent set* $\check{\mathbf{M}}$ in the obvious way.

Definition 3.1.5. Strictly \in -consistent ($\in_{\mathbf{s}}$ -consistent) set X it is a set such that

$$\forall x (x \in_{\mathbf{s}} X \vee x \notin_{\mathbf{s}} X).$$

Definition 3.1.6. Weakly \in -inconsistent (\in_w -inconsistent) set X it is a set such

$$\text{that } \forall x (x \in_w X \vee x \notin_w X).$$

Definition 3.1.7. Weakly \in -inconsistent with rank= n , $n \in \mathbb{N}$

($\in_{w,(n)}$ -inconsistent) set X it is a set such that:

$$\forall x (x \in_{w,(n)} X \vee x \notin_{w,(n)} X).$$

Definition 3.1.8. Strictly \in -inconsistent with rank= n , $n \in \mathbb{N}$

($\in_{w,[n]}$ -inconsistent) set X it is a set such that:

$$\forall x (x \in_{w,[n]} X \wedge x \notin_{w,[n]} X). \tag{3.1.5}$$

Definition 3.1.9. An strictly \in -consistent set $\check{\mathbf{A}} \subseteq_{\mathbf{s}} |\mathbf{M}|_{\text{inc}} =_{\mathbf{s}} |\check{\mathbf{M}}|$ is said

to be strictly consistent definable over $\check{\mathbf{M}}$ allowing parameters from $\check{\mathbf{B}}$ if there exists a formula $\varphi(n)$ with parameters from $\check{\mathbf{B}}$ and no free variables

other than n such that

$$\check{\mathbf{A}} = \{a \in_s |\check{\mathbf{M}}| : |\check{\mathbf{M}}| = \varphi(a), (\varphi(a))^{[0]} \vdash \mathbf{C}\}.$$

Here $\check{\mathbf{M}} \models \varphi(a)$ as it is usual means that $\check{\mathbf{M}}$ satisfies $\varphi(a)$, i.e., $\varphi(a)$ is true in $\check{\mathbf{M}}$ and \mathbf{C} is an any sentence of the inconsistent language $L_2^\#$.

Definition 3.1.10. The *intended model* for $L_2^\#$ is the model

$$\triangleq (\mathbb{N}^\#, \mathbf{P}(\mathbb{N}^\#), +, \times, 0_{\check{\mathbf{M}}} \triangleq \{0_s, 0_w, 0_{w,(i)}, 0_{w,[i]}\}, 1_{\check{\mathbf{M}}} \triangleq \{1_s, 1_w, 1_{w,(i)}, 1_{w,[i]}\}, <_{\check{\mathbf{M}}} \triangleq \{<_s, <_w, <_{w,(i)}, <_{w,[i]}\}, i \in \mathbb{N})$$

where $\mathbb{N}^\# \triangleq \mathbb{N}_{\text{inc}}$ is the set of paraconsistent natural numbers, $\mathbf{P}(\mathbb{N}^\#)$ is the set of all s -subsets of $\mathbb{N}^\#$, and $+$, \times , $0_{\check{\mathbf{M}}} \triangleq \{0_s, 0_w, 0_{w,(i)}, 0_{w,[i]}\}$, $1_{\check{\mathbf{M}}} \triangleq \{1_s, 1_w, 1_{w,(i)}, 1_{w,[i]}\}$, $<_{\check{\mathbf{M}}} \triangleq \{<_s, <_w, <_{w,(i)}, <_{w,[i]}\}$, $i \in \mathbb{N}$ are as below.

Definition 3.1.11. (second order paraconsistent arithmetic $\mathbf{Z}_2^\#$). The axioms of second order paraconsistent arithmetic $\mathbf{Z}_2^\#$ consist of the universal closures of the following $L_2^\#$ -formulas:

(i) basic axioms:

(i.1.) basic axioms of the first group:

- (i.1.1.) $n + 1_s \neq_s 0_s$,
- (i.1.2.) $n + 1_w \neq_s 0_w$,
- (i.1.3.) $n + 1_{w,(i)} \neq_s 0_{w,(i)}$, $i \in \mathbb{N}$,
- (i.1.4.) $n + 1_{w,[i]} \neq_s 0_{w,[i]}$, $i \in \mathbb{N}$,
- (i.1.5.) $1_s =_s 1_s$, $0_s =_s 0_s$,
- (i.1.6.) $1_w =_s 1_w$, $0_w =_s 0_w$,
- (i.1.7.) $1_{w,(i)} =_s 1_{w,(i)}$, $0_{w,(i)} =_s 0_{w,(i)}$, $i \in \mathbb{N}$,
- (i.1.8.) $1_{w,[i]} =_s 1_{w,[i]}$, $0_{w,[i]} =_s 0_{w,[i]}$, $i \in \mathbb{N}$;

(i.2.) basic axioms of the second group:

- (i.2.1.) $m + 1_s =_s n + 1_s \rightarrow m =_s n$,
- (i.2.2.) $m + 1_w =_w n + 1_w \rightarrow m =_w n$,
- (i.2.3.) $m + 1_{w,(i)} =_{w,(i)} n + 1_{w,(i)} \rightarrow m =_{w,(i)} n$, $i \in \mathbb{N}$,

(i.2.4.) $m + 1_{w,[i]} =_{w,[i]} n + 1_{w,[i]} \rightarrow m =_{w,[i]} n$, $i \in \mathbb{N}$;

(i.3.) basic axioms of the third group:

- (i.3.1.) $m + 0_s =_s m$,
- (i.3.2.) $m + 0_w =_w m$,
- (i.3.3.) $m + 0_{w,(i)} =_{w,(i)} m$, $i \in \mathbb{N}$,
- (i.3.4.) $m + 0_{w,[i]} =_{w,[i]} m$, $i \in \mathbb{N}$;

(i.4.) basic axioms of the fourth group:

- (i.4.1.) $m + (n + 1_s) =_s (m + n) + 1_s$,

(i.4.2.) $m + (n + 1_w) =_w (m + n) + 1_w$,

(i.4.3.) $m + (n + 1_{w,(i)}) =_{w,(i)} (m + n) + 1_{w,(i)}, i \in \mathbb{N}$,

(i.4.4.) $m + (n + 1_{w,[i]}) =_{w,[i]} (m + n) +_{w,[i]} 1_{w,[i]}, i \in \mathbb{N}$;

(i.5.) **basic axioms of the fifth group:**

(i.5.1.) $m \times 0_s =_s 0_s$,

(i.5.2.) $m \times 0_w =_w 0_w$,

(i.5.3.) $m \times 0_{w,(i)} =_{w,(i)} 0_{w,(i)}, i \in \mathbb{N}$,

(i.5.4.) $m \times 0_{w,[i]} =_{w,[i]} 0_{w,[i]}, i \in \mathbb{N}$;

(i.6.) **basic axioms of the sixth group:**

(i.6.1.) $m \times (n + 1_s) =_s (m \times n) + m$,

(i.6.2.) $m \times (n + 1_w) =_w (m \times n) + m$,

(i.6.3.) $m \times (n + 1_{w,(i)}) =_{w,(i)} (m \times n) + m, i \in \mathbb{N}$,

(i.6.4.) $m \times (n + 1_{w,[i]}) =_{w,[i]} (m \times n) + m, i \in \mathbb{N}$;

(i.7) **basic axioms of the seventh group:**

(i.7.1.) $\neg(m <_s 0_s)$,

(i.7.2.) $\neg(m <_w 0_w)$,

(i.7.3.) $\neg(m <_{w,(i)} 0_{w,(i)}), i \in \mathbb{N}$,

(i.7.4.) $\neg(m <_{w,[i]} 0_{w,[i]}), i \in \mathbb{N}$;

(i.8.) **basic axioms of the eighth group:**

(i.8.1.) $m <_s n + 1_s \leftrightarrow (m <_s n \vee m =_s n)$,

(i.8.2.) $m <_w n + 1_w \leftrightarrow (m <_w n \vee m =_w n)$,

(i.8.3.) $m <_{w,(i)} n + 1_{w,(i)} \leftrightarrow (m <_{w,(i)} n \vee m =_{w,(i)} n), i \in \mathbb{N}$,

(i.8.4.) $m <_{w,[i]} n + 1_{w,[i]} \leftrightarrow (m <_{w,[i]} n \vee m =_{w,[i]} n), i \in \mathbb{N}$;

Notation.3.1.1. $\forall n_{n \in \alpha} Z (n \in \alpha X) \iff \forall n [(n \in \alpha Z) \wedge (n \in \alpha X)]$.

(ii) **induction axioms:**

(ii.1.) **strictly consistent induction (s-induction) axiom**

$\exists Y_1 \forall X [((0_s \in_s X) \wedge \forall n_{n \in_s Y_1} [(n \in_s X \rightarrow n + 1_s \in_s X)] \rightarrow \forall n_{n \in_s Y_1} (n \in_s X)]$,

(ii.2.) **weakly inconsistent (w-inconsistent) induction axiom**

(*w*-induction axiom):

$\exists Y_2 \forall X [(0_w \in_w X \wedge \forall n_{n \in_w Y_2} [(n \in_w X \rightarrow n + 1_w \in_w X)] \rightarrow \forall n_{n \in_w Y_2} (n \in_w X)]$,

(ii.3.) **weakly inconsistent with rang=n, n ∈ N** ($\{w, (n)\}$ -inconsistent)

induction axiom ($\{w, (n)\}$ -induction axiom):

$\forall i_{(i \in \mathbb{N})} \exists Y_3^i \forall X \left[(0_{w,(i)} \in_{w,(i)} X \wedge \forall n_{n \in_{w,(i)} Y_3^i} [(n \in_{w,(i)} X \rightarrow n + 1_{w,(i)} \in_{w,(i)} X)] \rightarrow \forall n_{n \in_{w,(i)} Y_3^i} (n \in_{w,(i)} X) \right]$,

(ii.4.) **strictly inconsistent with rang=n, n ∈ N** ($\{w, [n]\}$ -inconsistent)

induction axiom ($\{w, [n]\}$ -induction axiom):

$\forall i_{(i \in \mathbb{N})} \exists Y_4^i \forall X \left[(0_{w,[i]} \in_{w,[i]} X \wedge \forall n_{n \in_{w,[i]} Y_4^i} [(n \in_{w,[i]} X \rightarrow n + 1_{w,[i]} \in_{w,[i]} X)] \rightarrow \forall n_{n \in_{w,[i]} Y_4^i} (n \in_{w,[i]} X) \right]$,

(ii.5.) $\{w, (\mathbb{N})\}$ -induction axiom:

$$\exists Y_5 \forall X [\forall i_{(i \in \mathbb{N})} [(0_{w,(i)} \in_{w,(i)} X \wedge \forall n_{n \in_{w,(i)} Y_5} [(n \in_{w,(i)} X \rightarrow n + 1_{w,(i)} \in_{w,(i)} X)])] \\ \rightarrow \forall j_{(j \in \mathbb{N})} \forall n_{n \in_{w,(j)} Y_5} (n \in_{w,(j)} X)],$$

(ii.6.) $\{w, [\mathbb{N}]\}$ -induction axiom:

$$\exists Y_6 \forall X [\forall i_{(i \in \mathbb{N})} [(0_{w,[i]} \in_{w,[i]} X \wedge \forall n_{n \in_{w,[i]} Y_6} [(n \in_{w,[i]} X \rightarrow n + 1_{w,[i]} \in_{w,[i]} X)])] \\ \rightarrow \forall j_{(j \in \mathbb{N})} \forall n_{n \in_{w,[j]} Y_6} (n \in_{w,[j]} X)],$$

(ii.7.) global paraconsistent induction axiom:

$$\exists Y_* \forall X [((0_s \in_s X) \wedge (0_w \in_s X) \wedge (\forall i_{(i \in \mathbb{N})} (0_{w,(i)} \in_s X)) \wedge (\forall i_{(i \in \mathbb{N})} (0_{w,[i]} \in_s X)) \wedge \\ \wedge \forall n_{n \in_s Y_*} (n \in_s X \rightarrow n + 1_s \in_s X) \wedge \forall n_{n \in_s Y_*} (n \in_s X \rightarrow n + 1_w \in_s X) \wedge \\ \wedge \{\forall i_{(i \in \mathbb{N})} \forall n_{n \in_s Y_*} (n \in_s X \rightarrow n + 1_{w,(i)} \in_s X)\} \wedge \{\forall i_{(i \in \mathbb{N})} \forall n_{n \in_s Y_*} (n \in_s X \rightarrow n + 1_{w,[i]} \in_s X)\} \\ \rightarrow \forall n_{n \in_s Y_*} (n \in_s X)].$$

$$Y_1 \triangleq \mathbb{N}_s^\# = \mathbb{N},$$

$$Y_2 \triangleq \mathbb{N}_w^\#,$$

$$Y_3^i \triangleq \mathbb{N}_{w,(i)}^\#,$$

Definition 3.1.12.

$$Y_4^i \triangleq \mathbb{N}_{w,[i]}^\#,$$

$$Y_5 \triangleq \mathbb{N}_{w,(\mathbb{N})}^\#,$$

$$Y_6 \triangleq \mathbb{N}_{w,[\mathbb{N}]}^\#,$$

$$Y_* \triangleq \mathbb{N}^\# \triangleq \mathbb{N}_{\mathbf{pc}}.$$

(iii) paraconsistent order axioms:

(iii.1.) weakly inconsistent order axiom (w-order axiom):

every nonempty w -subset $X \subseteq_w \mathbb{N}_w^\#$ has w -least element,i.e.

a least element relative to weakly inconsistent order $(\cdot <_w \cdot)$.

(iii.2.) weakly inconsistent with rank = n , $n \in \mathbb{N}$ order axiom

($\{w, (n)\}$ -order axiom):

every nonempty $\{w, (n)\}$ -subset $X \subseteq_{w,(n)} \mathbb{N}_{w,(n)}^\#$ has $\{w, (n)\}$ -least element,i.e.

a least element relative to weakly inconsistent order $(\cdot <_{w,(n)} \cdot)$.

(iii.3.) strictly inconsistent with rank = n , $n \in \mathbb{N}$ order axiom

($\{w, [n]\}$ -order axiom):

every nonempty $\{w, [n]\}$ -subset $X \subseteq_{w,[n]} \mathbb{N}_{w,(n)}^\#$ has $\{w, [n]\}$ -least element,i.e.

a least element relative to strictly inconsistent order $(\cdot <_{w,[n]} \cdot)$.

(iv) restricted comprehension scheme:

strictly consistent comprehension (s-comprehension) scheme:

(iv.1.) $\exists X \forall n_{n \in_s \mathbb{N}^\#} (n \in_s X \leftrightarrow \varphi(n)),$

where $\varphi(n)$ is any formula of $L_2^\#$ in which X does not occur freely.

(v.) non-restricted comprehension schemes:

(v.1.) weakly inconsistent comprehension scheme

(w-comprehension scheme):

$$\exists X \forall n_{n \in w \mathbb{N}_w^\#} (n \in_w X \leftrightarrow \varphi(n, X)),$$

(v.2.) weakly inconsistent with $\text{rank} =_n n \in \mathbb{N}$ comprehension scheme

($\{w, (n)\}$ -comprehension scheme):

$$\exists X \forall n (n \in_{w,(i)} X \leftrightarrow (\varphi(n, X))^{(i)}), i \in \mathbb{N},$$

(v.3.) strictly inconsistent with $\text{rank} =_n n \in \mathbb{N}$ comprehension scheme

($\{w, [n]\}$ -comprehension scheme) :

$$\exists X \forall n (n \in_{w,[i]} X \leftrightarrow (\varphi(n, X))^{[i]}), i \in \mathbb{N},$$

where $\varphi(n)$ is any formula of $L_2^\#$.

8 III.2.Paraconsistent Mathematics Within $\mathbf{Z}_2^\#$.

We now outline the development of certain portions of paraconsistent mathematics within $\mathbf{Z}_2^\#$.

Definition 3.2.1. If X and Y are set variables, we use:

(i) $X =_s Y$ and $X \subseteq_s Y$ as abbreviations for the formulas $\forall n (n \in_s X \leftrightarrow n \in_s Y)$

and $\forall n (n \in_s X \rightarrow n \in_s Y)$ respectively.

(ii) $X =_w Y$ and $X \subseteq_w Y$ as abbreviations for the formulas $\forall n (n \in_w X \leftrightarrow n \in_w Y)$

and $\forall n (n \in_w X \rightarrow n \in_w Y)$ respectively.

(iii) $X =_{w,(i)} Y$ and $X \subseteq_{w,(i)} Y, i \in \mathbb{N}$ as abbreviations for the formulas:

$$\forall n (n \in_{w,(i)} X \leftrightarrow n \in_{w,(i)} Y) \text{ and } \forall n (n \in_{w,(i)} X \rightarrow n \in_{w,(i)} Y)$$

respectively.

(iv) $X =_{w,[i]} Y$ and $X \subseteq_{w,[i]} Y, i \in \mathbb{N}$ as abbreviations for the formulas:

$$\forall n (n \in_{w,[i]} X \leftrightarrow n \in_{w,[i]} Y) \text{ and } \forall n (n \in_{w,[i]} X \rightarrow n \in_{w,[i]} Y)$$

respectively.

(v) a strictly \in -consistent (\in_s -consistent) set $X \subset_s \mathbb{N}^\#$ is defined to be a set

such that: $\forall n [(n \in_s X) \vee (n \notin_s X)]$

(vi) a weakly \in -inconsistent (\in_w -inconsistent) set $X \subset_w \mathbb{N}^\#$ is defined to be a

a

set such that: $\forall x (x \in_w X \vee x \notin_w X)$.

(vii) a weakly \in -inconsistent with rank $=_n n \in \mathbb{N}$

$(\in_{w,(n)}$ -inconsistent) set X it is a set $X \subset_{w,(n)} \mathbb{N}^\#$ such that:

$$\forall x (x \in_{w,(n)} X \vee x \notin_{w,(n)} X)$$

(viii) a strictly \in -inconsistent with rank $=_n n \in \mathbb{N}$

($\in_{w,[n]}$ -inconsistent) set X it is a set $X \subset_{w,[n]} \mathbb{N}^\#$ such that:

$$\forall x ((x \in_{w,[n]} X) \wedge (x \notin_{w,[n]} X))$$

Definition 3.2.2. Strictly consistent single element set:

$$\{x\}_s \triangleq \forall y [y \in_s \{x\}_s \longleftrightarrow$$

Weakly inconsistent single element set:

$$\{x\}_w \triangleq \forall y [y \in_w \{x\}_w \longleftrightarrow y =_w x].$$

Weakly inconsistent with rank= n , $n \in \mathbb{N}$ single element set:

$$\{x\}_{w,(n)} \triangleq \forall y [y \in_{w,(n)} \{x\}_{w,(n)} \longleftrightarrow$$

Strictly inconsistent with rank= n , $n \in \mathbb{N}$ single element set:

$$\{x\}_{w,[n]} \triangleq \forall y [y \in_{w,[n]} \{x\}_{w,[n]} \longleftrightarrow$$

Definition 3.2.3. Strictly consistent two-element set:

$$\{x, y\}_s \triangleq \forall z [z \in_s \{x, y\}_s \longleftrightarrow (z =_s x) \vee (z =_s y)]. \quad (3.2.5)$$

Weakly inconsistent two-element set:

$$\{x, y\}_w \triangleq \forall z [z \in_w \{x, y\}_w \longleftrightarrow (z =_w x) \vee (z =_w y)]$$

Weakly inconsistent with rank= n , $n \in \mathbb{N}$ two-element set:

$$\{x, y\}_{w,(n)} \triangleq \forall z [z \in_{w,(n)} \{x, y\}_{w,(n)} \longleftrightarrow$$

Strictly inconsistent with rank= n , $n \in \mathbb{N}$ two-element set:

$$\{x, y\}_{[n]} \triangleq \{x, y\}_{w,[n]} \triangleq \forall z [z \in_{w,[n]} \{x, y\}_{w,[n]} \longleftrightarrow (z =_{w,[n]} x) \vee (z =_{w,[n]} y)]. \quad (3.2.7)$$

Definition 3.2.4. Strictly consistent ordered pair:

$$(x, y)_s \triangleq \{\{x\}_s, \{x, y\}_s\}_s$$

Weakly inconsistent ordered pair:

$$(x, y)_w \triangleq \{\{x\}_w, \{x, y\}_w\}_w$$

Weakly inconsistent with rank= n , $n \in \mathbb{N}$ ordered pair:

$$(x, y)_{w,(n)} \triangleq \{\{x\}_{w,(n)}, \{x, y\}_{w,(n)}\}_{w,(n)}$$

Strictly inconsistent with rank= n , $n \in \mathbb{N}$ ordered pair:

$$(x, y)_{w,[n]} \triangleq \{\{x\}_{w,[n]}, \{x, y\}_{w,[n]}\}_{w,[n]} \quad (3.2.11)$$

Definition 3.2.5. The strictly consistent cartesian product of X and Y is:

$$X \times_s Y \triangleq \{(x, y)_s \mid (x \in_s X) \wedge (y \in_s Y)\}_s \quad (3.2.12)$$

The weakly inconsistent cartesian product of X and Y is:

$$X \times_w Y \triangleq \{(x, y)_w \mid (x \in_w X) \wedge (y \in_w Y)\}_w$$

The weakly inconsistent with rank= n , $n \in \mathbb{N}$ cartesian product

of X and Y is:

$$X \times_{w,(n)} Y \triangleq \{(x, y)_{w,(n)} \mid (x \in_{w,(n)} X) \wedge (y \in_{w,(n)} Y)\}_{w,(n)} \quad (3.2.13)$$

The strictly inconsistent cartesian product with rank= $n, n \in \mathbb{N}$ of X and Y

is:
$$X \times_{w,[n]} Y \triangleq \left\{ (x, y)_{w,[n]} \mid (x \in_{w,[n]} X)_{w,[n]} \wedge (y \in_{w,[n]} Y)_{w,[n]} \right\}_{w,[n]} \quad (3.2.15)$$

Definition 3.2.6. Within $\mathbf{Z}_2^\#$, we define $\mathbb{N}^\# \triangleq \mathbb{N}_{\mathbf{pc}}$ to be the unique set X such

that $\forall n \in_{\mathbf{s}} \mathbb{N}^\# (n \in_{\mathbf{s}} X)$.

Definition 3.2.7.(i) For $X, Y \subseteq_{\mathbf{s}} \mathbb{N}^\#$, a strictly consistent (\mathbf{s} -consistent)

function $f_{\mathbf{s}} : X \rightarrow_{\mathbf{s}} Y$ is defined to be a \in -consistent set $f \subseteq_{\mathbf{s}} X \times_{\mathbf{s}} Y$ such that

for all $m \in_{\mathbf{s}} X$ there is exactly one $n \in_{\mathbf{s}} Y$ such that $(m, n)_{\mathbf{s}} \in_{\mathbf{s}} f$. For $m \in_{\mathbf{s}} X$,

$f(m)$ is defined to be the \mathbf{s} -unique n such that $(m, n)_{\mathbf{s}} \in_{\mathbf{s}} f$.

The usual properties of such functions can be proved in $\mathbf{Z}_2^\#$.

(ii) For $X, Y \subseteq_w \mathbb{N}^\#$, a weakly inconsistent (w -inconsistent) function $f_w : X \rightarrow_w Y$

is defined to be a weakly \in -inconsistent set $f_w \subseteq_w X \times_w Y$ such that for all m such

that $(m \in_{\mathbf{s}} X) \vee (m \in_w X)$ there is exist w -unique $n \in_w Y$ such that $(m, n)_w \in_w f_w$.

For all m such that $(m \in_{\mathbf{s}} X) \vee (m \in_w X)$, $f_w(m)$ is defined to be the w -unique n

such that $(m, n)_w \in_w f_w$.

(iii) For $X, Y \subseteq_{w,(n)} \mathbb{N}_{w,(n)}^\#$, a weakly inconsistent with rank= $n, n \in \mathbb{N}$

($\{w, (n)\}$ -inconsistent) function $f_{w,(n)} : X \rightarrow_{w,(n)} Y$ is defined to be a weakly

(with rang= $n, n \in \mathbb{N}$) $\in_{w,(n)}$ -inconsistent set $f_{w,(n)} \subseteq_{w,(n)} X \times_{w,(n)} Y$ such that for all

m such that $(m \in_{\mathbf{s}} X) \vee (m \in_{w,(n)} X)$ there is exist $\{w, (n)\}$ -unique $n \in_{w,(n)} Y$ such

that $(m, n)_{w,(n)} \in_{w,(n)} f_{w,(n)}$. For $m \in_{w,(n)} X$, $f(m)$ is defined to be the $\{w, (n)\}$ -unique

n such that $(m, n)_{w,(n)} \in_{w,(n)} f_{w,(n)}$.

(iv) For $X, Y \subseteq_{w,[n]} \mathbb{N}^\#$, a strictly inconsistent with rank= $n, n \in \mathbb{N}$

($\{w, [n]\}$ -inconsistent) function $f_{w,[n]} : X \rightarrow_{w,[n]} Y$ is defined to be a strictly

inconsistent set $f_{w,[n]} \subseteq_{w,[n]} X \times_{w,[n]} Y$ such that for all $m \in_{w,[n]} X$ there is exist

$\{w, [n]\}$ -unique $n \in_{w,[n]} Y$ such that $(m, n)_{w,[n]} \in_{w,[n]} f$. For $m \in_{w,[n]} X$, $f(m)$ is

defined to be the $\{w, [n]\}$ -unique n such that $(m, n)_{w,[n]} \in_{w,[n]} f_{w,[n]}$.

Definition 3.2.8.(i) (the strictly consistent primitive recursion

(s-recursion)). This means that, given $f_s : X \rightarrow_s Y$ and $g_s : \mathbb{N}_s^\# \times_s X \times_s Y \rightarrow Y$,

there is a **s**-unique $h_s : \mathbb{N}_s^\# \times_s X \rightarrow_s Y$ defined by $h_s(0_s, m) = f_s(m)$,
 $h_s(n + 1_s, m) = g_s(n, m, h_s(n, m))$ for all $n \in_s \mathbb{N}_s^\#$ and $m \in_s X$.

The existence of h_s is proved by strictly consistent arithmetical comprehension,

and the **s**-uniqueness of h_s is proved by strictly consistent arithmetical

induction.

(ii) (the weakly consistent primitive recursion (w-recursion)). This means

that, given $f_w : X \rightarrow_w Y$ and $g_w : \mathbb{N} \times_w X \times_w Y \rightarrow_w Y$, there is a unique

$h_w : \mathbb{N}_w^\# \times_w X \rightarrow_w Y$ defined by $h_w(0_w, m) = f_w(m), h_w(n + 1_w, m) = g_w(n, m, h_w(n, m))$
for all $n \in_w \mathbb{N}_w^\#$ and $m \in_w X$.

The existence of h_w is proved by weakly consistent arithmetical comprehension,

and the **w**-uniqueness of h_w is proved by weakly consistent arithmetical

induction.

(iii) (the weakly consistent with rank= i , $i \in \mathbb{N}$ primitive recursion

($\{w, (n)\}$ -recursion)). This means that, given $f_{w,(n)} : X \rightarrow_{w,(n)} Y$ and
 $g_{w,(n)} : \mathbb{N}_{w,(n)}^\# \times_{w,(n)} X \times_{w,(n)} Y \rightarrow_{w,(n)} Y$, there is a unique $h_s : \mathbb{N}_{w,(n)}^\# \times_{w,(n)} X \rightarrow_{w,(n)} Y$
defined by $h_{w,(n)}(0_{w,(n)}, m) = f_{w,(n)}(m), h_{w,(n)}(n + 1_{w,(n)}, m) = g_{w,(n)}(n, m, h_{w,(n)}(n, m))$
for all

$n \in_{w,(n)} \mathbb{N}_{w,(n)}^\#$ and $m \in_{w,(n)} X$.

The existence of $h_{w,(n)}$ is proved by weakly consistent with rank= n , $n \in \mathbb{N}$

arithmetical comprehension and the uniqueness of $h_{w,(n)}$ is proved by weakly consistent with rank= n , $n \in \mathbb{N}$ arithmetical induction.

(iv) (the strictly inconsistent with rank= i , $i \in \omega$ primitive recursion

($\{w, [n]\}$ -recursion)). This means that, given $f_{w,[n]} : X \rightarrow_{w,[n]} Y$ and
 $g_{w,[n]} : \mathbb{N} \times_{w,[n]} X \times_{w,[n]} Y \rightarrow_{w,[n]} Y$, there is a unique $h_s : \mathbb{N}_{w,[n]}^\# \times_{w,[n]} X \rightarrow_{w,[n]} Y$
defined by $h_{w,[n]}(0_{w,[n]}, m) = f_{w,[n]}(m), h_{w,[n]}(n + 1_{w,[n]}, m) = g_{w,[n]}(n, m, h_{w,[n]}(n, m))$

for all $n \in_{w,[n]} \mathbb{N}_{w,[n]}^\#$ and $m \in_{w,[n]} X$.

The existence of $h_{w,[n]}$ is proved by strictly inconsistent with rank= n , $n \in \mathbb{N}$

arithmetical comprehension and the uniqueness of $h_{w,[n]}$ is proved by strictly inconsistent with rank= n , $n \in \omega$ arithmetical induction.

(v) (the global paraconsistent primitive recursion). This means that,

given $f_{\text{gl}}^s : X \rightarrow_s Y$, $f_{\text{gl}}^w : X \rightarrow_w Y$, $f_{\text{gl}}^{w.(i)} : X \rightarrow_s Y$, $f_{\text{gl}}^{w,[i]} : X \rightarrow_s Y$, $i \in \mathbb{N}$ and

given $g_{\text{gl}}^s : \mathbb{N}^\# \times_s X \times_s Y \rightarrow_s Y$, $g_{\text{gl}}^w : \mathbb{N}^\# \times_s X \times_s Y \rightarrow_s Y$,
 $g_{\text{gl}}^{w,(i)} : \mathbb{N}^\# \times_s X \times_s Y \rightarrow_s Y$, $g_{\text{gl}}^{w,[i]} : \mathbb{N}^\# \times_s X \times_s Y \rightarrow_s Y$, $i \in \mathbb{N}$ there is a weakly

unique $h_{\text{gl}} : \mathbb{N}^\# \times_s X \rightarrow_s Y$ defined by:
 $h_{\text{gl}}^s(0_s, m) =_s f_{\text{gl}}^s(m)$, $h_{\text{gl}}^w(0_w, m) =_s f_{\text{gl}}^w(m)$,
 $h_{\text{gl}}^{w,(i)}(0_{w,(i)}, m) =_s f_{\text{gl}}^{w,(i)}(m)$,
 $h_{\text{gl}}^{w,[i]}(0_{w,[i]}, m) =_s f_{\text{gl}}^{w,[i]}(m)$, $i \in \mathbb{N}$,
 $h_{\text{gl}}^s(n + 1_s, m) =_s g_{\text{gl}}^s(n, m, h_s(n, m))$,
 $h_{\text{gl}}^s(n + 1_w, m) =_s g_{\text{gl}}^s(n, m, h_{\text{gl}}(n, m))$,
 $h_{\text{gl}}^{w,(i)}(n + 1_{w,(i)}, m) =_s g_{\text{gl}}^{w,(i)}(n, m, h_{\text{gl}}(n, m))$,
 $h_{\text{gl}}^{w,[i]}(n + 1_{w,[i]}, m) =_s g_{\text{gl}}^{w,[i]}(n, m, h_s(n, m))$,
 $i \in \mathbb{N}$, for all $n \in_s \mathbb{N}^\#$ and $m \in_s X$.

The existence of h_s is proved by strictly consistent arithmetical comprehension and the uniqueness of h_s is proved by global paraconsistent arithmetical induction.

Remark 3.1.1. In particular, we have:

(i) the strictly consistent exponential function $\exp(m, n)_s =_s (m)_s^n$,

defined by $(m)_s^{0_s} =_s 1_s$,

$(m)_s^{n+1_s} =_s (m)_s^n \times_s m$ for all $m, n \in_s \mathbb{N}_s^\#$.

(ii) the weakly inconsistent exponential function $\exp_w(m, n) =_w (m)_w^n$,

defined by $(m)_w^{0_w} =_w 1_w$, $(m)_w^{n+1_w} =_w (m)_w^n \times_w m$ for all $m, n \in_s \mathbb{N}_w^\#$.

(iii) the weakly inconsistent with rank= i , $i \in \mathbb{N}$ exponential function $\exp_{w,(i)}(m, n) =_w$

$(m)_{w,(i)}^n$, defined by $(m)_{w,(i)}^{0_{w,(i)}} =_s 1_{w,(i)}$,

$(m)_{w,(i)}^{n+1_{w,(i)}} =_{w,(i)} (m)_{w,(i)}^n \times_w m$ for all $m, n \in_{w,(i)} \mathbb{N}_{w,(i)}^\#$.

(iv) the strictly inconsistent with rank= i , $i \in \mathbb{N}$ exponential function $\exp_{w,[i]}(m, n) =_{w,[i]}$

$(m)_{w,[i]}^n$, defined by $(m)_{w,[i]}^{0_{w,[i]}} =_{w,[i]} 1_{w,[i]}$,

$(m)_{w,[i]}^{n+1_{w,[i]}} =_{w,[i]} (m)_{w,[i]}^n \times m$ for all $m, n \in_{w,[i]} \mathbb{N}_{w,[i]}^\#$.

The usual properties of the exponential function can be proved in $\mathbf{Z}_2^\#$.

(v) the global paraconsistent exponential function $\exp_{\text{gl}}(m, n) \triangleq (m)_{\text{gl}}^n$,

defined by: $(m)_{\text{gl}}^{0_s} =_s 1_s$, $(m)_{\text{gl}}^{n+1_s} =_s (m)_{\text{gl}}^n \times m$ for all $m, n \in_s \mathbb{N}_s^\#$,

$(m)_{\text{gl}}^{0_w} =_w 1_w$, $(m)_{\text{gl}}^{n+1_w} =_w (m)_{\text{gl}}^n \times m$ for all $m, n \in_s \mathbb{N}_w^\#$,

$(m)_{\text{gl}}^{0_{w,(n)}} =_{w,(n)} 1_{w,(n)}$, $(m)_{\text{gl}}^{n+1_{w,(n)}} =_{w,(n)} (m)_{\text{gl}}^n \times m$ for all $m, n \in_s \mathbb{N}_{w,(n)}^\#$,

$(m)_{\text{gl}}^{0_{w,[n]}} =_{w,[n]} 1_{w,[n]}$, $(m)_{\text{gl}}^{n+1_{w,[n]}} =_{w,[n]} (m)_{\text{gl}}^n \times m$ for all $m, n \in_s \mathbb{N}_{w,[n]}^\#$.

The usual properties of the exponential function can be proved in $\mathbf{Z}_2^\#$.

Within $\mathbf{Z}_2^\#$, we define a numerical strictly consistent pairing function $\pi_s(m, n)$ by $\pi_s(m, n) =_s (m + n)^2 + m$. Within $\mathbf{Z}_2^\#$ we can prove that, for

all $m, n, i, j \in_s \mathbb{N}^\#$, $\pi_s(m, n) =_s \pi_s(i, j)$ if and only if $m =_s i$ and $n =_s j$.

Moreover, using strictly consistent arithmetical comprehension, we can prove that for all sets $X, Y \subseteq_s \mathbb{N}^\#$, there exists a set $\pi_s(X \times Y) \subseteq_s \mathbb{N}^\#$

consisting of all $\pi_s(m, n)$ such that $m \in_s X$ and $n \in_s Y$. In particular we

have $\pi_s(\mathbb{N}^\# \times \mathbb{N}^\#) \subseteq_s \mathbb{N}^\#$.

The paraconsistent natural number system is essentially already given to us by the language $L_2^\#$ and axioms of $\mathbf{Z}_2^\#$. Thus, within $\mathbf{Z}_2^\#$, a consistent

and inconsistent natural number is defined to be an element of $\mathbb{N}^\#$, and

the *paraconsistent natural number system* is defined to be the

$$\begin{aligned} \mathbb{N}^\#, +_{\mathbb{N}^\#}, \times_{\mathbb{N}^\#}, 0_{\mathbb{N}^\#}, 1_{\mathbb{N}^\#}, & <_{\mathbb{N}^\#}, =_{\mathbb{N}^\#}, \\ 0_{\mathbb{N}^\#} & \triangleq \{0_s, 0_w, 0_{w,(i)}, 0_{w,[i]}\}, \\ 1_{\mathbb{N}^\#} & \triangleq \{1_s, 1_w, 1_{w,(i)}, 1_{w,[i]}\}, \\ <_{\mathbb{N}^\#} & \triangleq \{<_s, <_w, <_{w,(i)}, <_{w,[i]}\} \\ =_{\mathbb{N}^\#} & \triangleq \{=_s, =_w, =_{w,(i)}, =_{w,[i]}\} \end{aligned}$$

where $+_{\mathbb{N}^\#} : \mathbb{N}^\# \times_s \mathbb{N}^\# \rightarrow_s \mathbb{N}^\#$ is defined by $m +_{\mathbb{N}^\#} n =_s m + n$, etc.

Thus for instance $+_{\mathbb{N}^\#}$ is the set of triples $((m, n)_s, k)_s \in_s (\mathbb{N}^\# \times_s \mathbb{N}^\#) \times_s \mathbb{N}^\#$ such that $m + n =_s k$. The existence of this set follows from the strictly

consistent arithmetical comprehension.

In a standard manner, we can define within $\mathbf{Z}_2^\#$ the set $\mathbb{Z}^\# \triangleq \mathbb{Z}_{\text{pc}}$ of the all

consistent and inconsistent integers and the set of the all consistent and inconsistent rational numbers: $\mathbb{Q}^\# \triangleq \mathbb{Q}_{\text{pc}}$.

Definition 3.2.9. (paraconsistent rational numbers $\mathbb{Q}^\#$) Let

be the set of positive consistent integers, and let

(a) $\equiv_{\mathbb{Q}^\#, s}$ be the strictly equivalence relation on $\mathbb{Z}^\# \times_s (\mathbb{Z}^\#)^+$ defined by $(a, b)_s \equiv_{\mathbb{Q}^\#} (c, d)_s$ if and only if $a \times_{\mathbb{Z}^\#} d =_s b \times_{\mathbb{Z}^\#} c$. Then $\mathbb{Q}^\#$ is defined to be the set of all $(a, b)_s \in \mathbb{Z}^\# \times_s (\mathbb{Z}^\#)^+$ such

that $(a, b)_{\mathbf{s}}$ is the $<_{\mathbf{s}}^{\mathbb{N}^{\#}}$ -minimum element of its $\equiv_{\mathbb{Q}^{\#}}$ -equivalence class.

Operations $+_{\mathbb{Q}^{\#}}, -_{\mathbb{Q}^{\#}}, \times_{\mathbb{Q}^{\#}}$ on $\mathbb{Q}^{\#}$ are defined by:

$$(a, b)_{\mathbf{s}} +_{\mathbb{Q}^{\#}} (c, d)_{\mathbf{s}} \equiv_{\mathbb{Q}^{\#}} (a \times_{\mathbb{Z}^{\#}} d +_{\mathbb{Z}^{\#}} b \times_{\mathbb{Z}^{\#}} c, b \times_{\mathbb{Z}^{\#}} d)_{\mathbf{s}},$$

$$-_{\mathbb{Q}^{\#}}(a, b)_{\mathbf{s}} \equiv_{\mathbb{Q}^{\#}} (-_{\mathbb{Z}^{\#}}a, b)_{\mathbf{s}}, \text{ and}$$

$$(a, b)_{\mathbf{s}} \times_{\mathbb{Q}^{\#}} (c, d)_{\mathbf{s}} \equiv Q(a \times_{\mathbb{Z}^{\#}} c, b \times_{\mathbb{Z}^{\#}} d)_{\mathbf{s}}. \text{ We let } 0_{\mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (0_{\mathbb{Z}^{\#}}, 1_{\mathbb{Z}^{\#}})_{\mathbf{s}}, \text{i.e.}$$

$$0_{\mathbf{s}, \mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (0_{\mathbf{s}, \mathbb{Z}^{\#}}, 1_{\mathbf{s}, \mathbb{Z}^{\#}})_{\mathbf{s}}, 0_{w, \mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (0_{w, \mathbb{Z}^{\#}}, 1_{w, \mathbb{Z}^{\#}})_{\mathbf{s}}, 0_{w, (n), \mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (0_{w, (n), \mathbb{Z}^{\#}}, 1_{w, (n), \mathbb{Z}^{\#}})_{\mathbf{s}},$$

$$0_{w, [n], \mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (0_{w, [n], \mathbb{Z}^{\#}}, 1_{w, [n], \mathbb{Z}^{\#}})_{\mathbf{s}},$$

and

$$1_{\mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (1_{\mathbb{Z}^{\#}}, 1_{\mathbb{Z}^{\#}})_{\mathbf{s}}, \text{i.e. } 1_{\mathbf{s}, \mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (1_{\mathbf{s}, \mathbb{Z}^{\#}}, 1_{\mathbf{s}, \mathbb{Z}^{\#}})_{\mathbf{s}}, 1_{w, \mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (1_{w, \mathbb{Z}^{\#}}, 1_{w, \mathbb{Z}^{\#}})_{\mathbf{s}},$$

$$1_{w, (n), \mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (1_{w, (n), \mathbb{Z}^{\#}}, 1_{w, (n), \mathbb{Z}^{\#}})_{\mathbf{s}}, 1_{w, [n], \mathbb{Q}^{\#}} \equiv_{\mathbb{Q}^{\#}} (1_{w, [n], \mathbb{Z}^{\#}}, 1_{w, [n], \mathbb{Z}^{\#}})_{\mathbf{s}},$$

and we define a binary relations $<_{\mathbb{Q}^{\#}} \triangleq \{<_{\mathbf{s}}, <_w, <_{w, (i)}, <_{w, [i]}\}$ on $\mathbb{Q}^{\#}$ by

letting $(a, b)_{\mathbf{s}} <_{\mathbb{Q}^{\#}} (c, d)_{\mathbf{s}}$ if

and only if $a \times_{\mathbb{Z}^{\#}} d <_{\mathbb{Z}^{\#}} b \times_{\mathbb{Z}^{\#}} c$. Finally $=_{\mathbb{Q}^{\#}} \triangleq \{=_{\mathbf{s}}, =_w, =_{w, (i)}, =_{w, [i]}\}$ is the

identity relations on $\mathbb{Q}^{\#}$. We can then prove within $\mathbf{Z}_2^{\#}$ that the paraconsistent

rational number system

$$\begin{aligned} & \{\mathbb{Q}^{\#}, +_{\mathbb{Q}^{\#}}, -_{\mathbb{Q}^{\#}}, \times_{\mathbb{Q}^{\#}}, 0_{\mathbb{Q}^{\#}}, 1_{\mathbb{Q}^{\#}}, <_{\mathbb{Q}^{\#}}, =_{\mathbb{Q}^{\#}}\}, \\ & 0_{\mathbb{Q}^{\#}} \triangleq \{0_{\mathbf{s}, \mathbb{Q}^{\#}}, 0_{w, \mathbb{Q}^{\#}}, 0_{w, (n), \mathbb{Q}^{\#}}, 0_{w, [n], \mathbb{Q}^{\#}}\}, \\ & 1_{\mathbb{Q}^{\#}} \triangleq \{1_{\mathbf{s}, \mathbb{Q}^{\#}}, 1_{w, \mathbb{Q}^{\#}}, 1_{w, (n), \mathbb{Q}^{\#}}, 1_{w, [n], \mathbb{Q}^{\#}}\}, \\ & <_{\mathbb{Q}^{\#}} \triangleq \{<_{\mathbf{s}, \mathbb{Q}^{\#}}, <_w, <_{w, (i), \mathbb{Q}^{\#}}, <_{w, [i], \mathbb{Q}^{\#}}\}, \\ & =_{\mathbb{Q}^{\#}} \triangleq \{=_{\mathbf{s}, \mathbb{Q}^{\#}}, =_w, =_{w, (i), \mathbb{Q}^{\#}}, =_{w, [i], \mathbb{Q}^{\#}}\}, \end{aligned} \tag{3.2.18}$$

has the usual properties of an paraordered field, etc.

We make the usual identifications whereby $\mathbb{N}^{\#}$ is regarded as a

subset of $\mathbb{Z}^{\#}$ and $\mathbb{Z}^{\#}$ is regarded as a subset of $\mathbb{Q}^{\#}$. Namely $m \in_{\mathbf{s}} \mathbb{N}^{\#}$ is

identified with $(m, 0)_{\mathbf{s}} \in_{\mathbf{s}} \mathbb{Z}^{\#}$, and $a \in_{\mathbf{s}} \mathbb{Z}^{\#}$ is identified with $(a, 1_{\mathbb{Z}^{\#}})_{\mathbf{s}} \in_{\mathbf{s}} \mathbb{Q}^{\#}$.

We use $+$ ambiguously to denote $+_{\mathbb{N}^{\#}}, +_{\mathbb{Z}^{\#}}$, or $+_{\mathbb{Q}^{\#}}$ and

similarly for $-$, \times , 0 , 1 , $<$. For $q, r \in_{\mathbf{s}} \mathbb{Q}^{\#}$ we write $q - r = q + (-r)$, and if $r \neq_{\mathbb{Q}^{\#}} 0$, $q/r =_{\mathbb{Q}^{\#}}$ the unique $q' \in_{\mathbf{s}} \mathbb{Q}^{\#}$ such that $q = q' \times r$. The function $\mathbf{s}\text{-exp}(q, a) = \mathbf{s}\text{-}q^a$

for $q \in_{\mathbf{s}} \mathbb{Q}^{\#} \setminus \{0_{\mathbb{Q}^{\#}}\}$ and $a \in_{\mathbf{s}} \mathbb{Z}^{\#}$ is obtained by primitive recursion in the obvious

way.

Definition 3.2.10. The absolute value functions:

(i) the strictly consistent absolute value function:

$|\cdot|_{\mathbf{s}} : \mathbb{Q}^{\#} \rightarrow_{\mathbf{s}} \mathbb{Q}^{\#}$ is defined by $|q|_{\mathbf{s}} =_{\mathbf{s}} q$ if $0 \leq_{\mathbf{s}} q$, and by $|q|_{\mathbf{s}} =_{\mathbf{s}} -q$ otherwise,

(ii) the weakly inconsistent absolute value function:

$|\cdot|_w : \mathbb{Q}^\# \rightarrow_s \mathbb{Q}^\#$ is defined by $|q|_w =_s q$ if $0 \leq_w q$, and by $|q|_w =_s -q$ otherwise,

(iii) the weakly inconsistent with $\text{rang}=n, n \in \omega$ absolute value function:

$|\cdot|_{w,(n)} : \mathbb{Q}^\# \rightarrow_s \mathbb{Q}^\#$ is defined by $|q|_{w,(n)} =_s q$ if $0 \leq_{w,(n)} q$, and by $|q|_w =_s -q$ otherwise,

(iv) the strictly inconsistent with $\text{rang}=n, n \in \omega$ absolute value function:

$|\cdot|_{w,[n]} : \mathbb{Q}^\# \rightarrow_s \mathbb{Q}^\#$ is defined by $|q|_{w,[n]} =_s q$ if $0 \leq_{w,[n]} q$, and by $|q|_{w,[n]} =_s -q$ otherwise.

Definition 3.2.11.(a) A strictly consistent sequence of paraconsistent rational

numbers is defined to be a strictly consistent function $f_s : \mathbb{N}^\# \rightarrow_s \mathbb{Q}^\#$.

We denote such a strictly consistent sequence as $\langle q_n : n \in_s \mathbb{N}^\# \rangle_s$, or

simply $\langle q_n \rangle_s$, where $q_n =_s f_s(n)$.

Definition 3.2.12. A double strictly consistent sequence of paraconsistent

numbers is a consistent function $f_s : \mathbb{N}^\# \times_s \mathbb{N}^\# \rightarrow \mathbb{Q}^\#$, denoted

$\langle q_{mn} : m, n \in_s \mathbb{N}^\# \rangle_s$ or simply $\langle q_{mn} \rangle_s$, where $q_{mn} =_s f_s(m, n)$.

Definition 3.2.13.(paraconsistent real numbers).(a) Within $\mathbb{Z}_2^\#$, a strictly

consistent real number is defined to be a Cauchy strictly consistent sequence

of paraconsistent rational numbers, i.e., as a strictly consistent sequence of

paraconsistent rational numbers $x =_s \langle q_n : n \in_s \mathbb{N}^\# \rangle_s$ such that:

$\forall \varepsilon (\varepsilon \in_s \mathbb{Q}^\#)$

(b) Within $\mathbb{Z}_2^\#$, a weakly inconsistent real number is defined to be a Cauchy

strictly consistent sequence of paraconsistent rational numbers, i.e., as a

strictly consistent sequence of paraconsistent rational numbers

$x =_s \langle q_n : n \in_s \mathbb{N}^\# \rangle_s$ such that:

$\forall \varepsilon (\varepsilon \in_s \mathbb{Q}^\#) (0_w <_w \varepsilon \rightarrow \exists m \forall n (m <_w n$

(c) Within $\mathbb{Z}_2^\#$, a weakly inconsistent with $\text{rang}=n, n \in \omega$ real number

is

defined to be a Cauchy strictly consistent sequence of paraconsistent rational numbers, i.e., as a strictly consistent sequence of paraconsistent

rational numbers $x =_s \langle q_n : n \in_s \mathbb{N}^\# \rangle_s$ such that:

$\forall \varepsilon (\varepsilon \in_s \mathbb{Q}^\#) (0_{w,(n)} <_{w,(n)} \varepsilon \rightarrow \exists m \forall n (m <_{w,(n)} n$

(d) Within $\mathbb{Z}_2^\#$, a strictly inconsistent with $\text{rang}=n, n \in \omega$ real number is defined

to be a Cauchy strictly consistent sequence of paraconsistent rational numbers, i.e., as a strictly consistent sequence of paraconsistent rational

numbers $x =_s \langle q_n : n \in_s \mathbb{N}^\# \rangle_s$ such that:

$$\forall \varepsilon (\varepsilon \in_s \mathbb{Q}^\#) (0_{w,[n]} <_{w,[n]} \varepsilon \rightarrow \exists m \forall n (m$$

Definition 3.2.14.(a) If $x =_s q_n$ and $y =_s q'_n$ are strictly consistent real

numbers, we write $x =_{s,\mathbb{R}^\#} y$ to mean that $\mathbf{s}\text{-}\lim_n |q_n - q'_n|_s =_s 0_s$, i.e.,

$$\forall \varepsilon (\varepsilon \in_s \mathbb{Q}^\#) (0_s <_s \varepsilon \rightarrow \exists m \forall n (m <_s n \rightarrow |q_n - q'_n|_s <_s \varepsilon)), \quad (3.2.23)$$

and we write $x <_{s,\mathbb{R}^\#} y$ to mean that

$$\exists \varepsilon (0 <_{s,\mathbb{R}^\#} \varepsilon \wedge \exists m \forall n (m <_s n \rightarrow q_n + \varepsilon <_s q'_n)).$$

(b) If $x =_w q_n$ and $y =_w q'_n$ are weakly consistent real numbers, we

write $x =_{w,\mathbb{R}^\#} y$ to mean that $w\text{-}\lim_n |q_n - q'_n|_w =_w 0_w$, i.e.,

$$\forall \varepsilon (\varepsilon \in_s \mathbb{Q}^\#) (0_w <_w \varepsilon \rightarrow \exists m \forall n (m <_w n \rightarrow |q_n - q'_n|_w <_w \varepsilon)).$$

and we write $x <_{w,\mathbb{R}^\#} y$ to mean that

$$\exists \varepsilon (0 <_w \varepsilon \wedge \exists m \forall n (m <_w n \rightarrow q_n + \varepsilon <_w q'_n)).$$

(c) If $x =_{w,(n)} q_n$ and $y =_{w,(n)} q'_n$ are weakly inconsistent with $\text{rang}=n, n \in \omega$ real numbers, we write $x =_{w,(n),\mathbb{R}^\#} y$ to mean

that $\{w, (n)\} \text{-}\lim_n |q_n - q'_n|_{w,(n)} =_{w,(n)} 0_{w,(n)}$, i.e.,

$$\forall \varepsilon (\varepsilon \in_s \mathbb{Q}^\#) (0_{w,(n)} <_{w,(n)} \varepsilon \rightarrow \exists m \forall n (m <_s n \rightarrow |q_n - q'_n|_{w,(n)} <_s \varepsilon)).$$

and we write $x <_{w,(n),\mathbb{R}^\#} y$ to mean that

$$\exists \varepsilon (0 <_{w,(n)} \varepsilon \wedge \exists m \forall n (m <_s n \rightarrow q_n + \varepsilon <_{w,(n)} q'_n)).$$

(d) If $x =_s q_n$ and $y =_s q'_n$ are strictly inconsistent with $\text{rang}=n, n \in \omega$ real numbers, we write $x =_{\mathbb{R}^\#, w, [n]} y$ to mean

that $w, [n] \text{-}\lim_n |q_n - q'_n|_{w, [n]} =_{w, [n]} 0_{w, [n]}$, i.e.,

$$\forall \varepsilon (\varepsilon \in_s \mathbb{Q}^\#) (0_{w, [n]} <_{w, [n]} \varepsilon \rightarrow \exists m \forall n (m <_s n \rightarrow |q_n - q'_n|_{w, [n]} <_s \varepsilon)).$$

and we write $x <_{w, [n], \mathbb{R}^\#} y$ to mean that

$$\exists \varepsilon (0 <_{w, [n]} \varepsilon \wedge \exists m \forall n (m <_s n \rightarrow q_n + \varepsilon <_{w, [n]} q'_n)).$$

Also $x +_{\mathbb{R}^\#} y = \langle q_n + q'_n \rangle_s$, $x \times_{\mathbb{R}^\#} y =_s \langle q_n \times q'_n \rangle_s$, $-_{\mathbb{R}^\#} x =_s \langle -q_n \rangle_s$,

$0_{s, \mathbb{R}^\#} = \langle 0_s \rangle_s$, $1_{s, \mathbb{R}^\#} = \langle 1_s \rangle_s$, etc.

We use $\mathbb{R}^\#$ to denote the set of all *paraconsistent real numbers*.

Thus $x \in_s \mathbb{R}^\#$ means that x is a *consistent real number*. (Formally, we cannot speak of the set $\mathbb{R}^\#$ within the language of second order arithmetic, since it is a set of sets.)

We shall usually omit the subscript $\mathbb{R}^\#$ in $+_{\mathbb{R}^\#}$, $-_{\mathbb{R}^\#}$, $\times_{\mathbb{R}^\#}$, $0_{s, \mathbb{R}^\#}$, $0_{w, \mathbb{R}^\#}$, $0_{w, (n), \mathbb{R}^\#}$, $0_{w, [n], \mathbb{R}^\#}$, $1_{s, \mathbb{R}^\#}$, $1_{w, \mathbb{R}^\#}$, $1_{w, (n), \mathbb{R}^\#}$, $1_{w, [n], \mathbb{R}^\#}$, $<_{s, \mathbb{R}^\#}$, $<_{w, \mathbb{R}^\#}$, $<_{w, (n), \mathbb{R}^\#}$, $<_{w, [n], \mathbb{R}^\#}$,

$=_{s, \mathbb{R}^\#}$, $=_{w, \mathbb{R}^\#}$, $=_{w, (n), \mathbb{R}^\#}$, $=_{w, [n], \mathbb{R}^\#}$.

Thus the *consistent real number system* consists of $\mathbb{R}^\#, +, -, \times, 0_s, 0_w,$

$0_{w, (n)}, 0_{w, [n]}, 1_s, 1_w, 1_{w, (n)}, 1_{w, [n]}, <_s, <_w, <_{w, (n)}, <_{w, [n]}, =_s, =_w, =_{w, (n)}, =_{w, [n]}$.

We shall sometimes identify a paraconsistent rational number $q \in_s \mathbb{Q}^\#$

with the corresponding paraconsistent real number $x_q =_s \langle q \rangle_s \in_s \mathbb{R}^\#$.

Within $\mathbf{Z}_2^\#$ one can prove that the real number system has the usual

properties of an *paraconsistent Archimedean paraordered field*, etc. The

complex paraconsistent numbers can be introduced as usual as pairs of paraconsistent real numbers.

Within $\mathbf{Z}_2^\#$, it is straightforward to carry out the proofs of all the basic results

in real and complex *paraconsistent linear* and *paraconsistent polynomial*

algebra. For example, the paraconsistent analog of the fundamental theorem

of algebra can be proved in $\mathbf{Z}_2^\#$.

Definition 3.2.15. A strictly consistent sequence of paraconsistent real

numbers is defined to be a double strictly consistent sequence of rational

numbers $\langle q_{mn} : m, n \in_s \mathbb{N}^\# \rangle_s$ such that for each m , $\langle q_{mn} : n \in_s \mathbb{N}^\# \rangle_s$ is a

consistent real number.

Such a strictly consistent sequence of paraconsistent real numbers is denoted $\langle x_m : m \in_s \mathbb{N}^\# \rangle_s$, where $x_m =_s \langle q_{mn} : n \in_s \mathbb{N}^\# \rangle_s$. Within $\mathbf{Z}_2^\#$ we can prove that every

bounded (in *paraconsistent sense*) strictly consistent sequence of paraconsistent real numbers has a *paraconsistent least upper bound*. This is

a

very useful *paracompleteness* property of the paraconsistent real number

system.

For instance, it implies that an infinite series of positive terms is

paraconvergent if and only if the finite partial sums are *parabounded*.

We now turn of certain portions of paraconsistent abstract algebra within $\mathbf{Z}_2^\#$.

Because of the restriction to the language $L_2^\#$ of second order paraconsistent arithmetic, we cannot expect to obtain a good general theory of arbitrary

(countable and uncountable) paraconsistent algebraic structures. However, we

can develop paraconsistent countable algebra, i.e., the theory of countable paraconsistent algebraic structures, within $\mathbf{Z}_2^\#$.

Definition 3.2.16. A countable paraconsistent commutative ring is defined

within $\mathbf{Z}_2^\#$ to be a paraconsistent structure $\mathbf{R}_{\text{inc}}, +_{\mathbf{R}_{\text{inc}}}, -_{\mathbf{R}_{\text{inc}}}, \times_{\mathbf{R}_{\text{inc}}}$,

$0_{s, \mathbf{R}_{\text{inc}}}, 0_{w, \mathbf{R}_{\text{inc}}}, 0_{w, (n), \mathbf{R}_{\text{inc}}}, 0_{w, [n], \mathbf{R}_{\text{inc}}}, 1_{s, \mathbf{R}_{\text{inc}}}, 1_{w, \mathbf{R}_{\text{inc}}}, 1_{w, (n), \mathbf{R}_{\text{inc}}}, 1_{w, [n], \mathbf{R}_{\text{inc}}}$ where $\mathbf{R}_{\text{inc}} \subseteq_s \mathbb{N}^\#$, $+_{\mathbf{R}_{\text{inc}}} : \mathbf{R}_{\text{inc}} \times_s \mathbf{R}_{\text{inc}} \rightarrow_s \mathbf{R}_{\text{inc}}$, etc., and the usual commutative

paraconsistent ring axioms are assumed.

(We include $0_s \neq_s 1_s, 0_w \neq_s 1_w, 0_{w,(n)} \neq_s 1_{w,(n)}, 0_{w,[n]} \neq_s 1_{w,[n]}$, among those axioms.) The subscript \mathbf{R}_{inc} is usually omitted.

An strictly consistent ideal in \mathbf{R}_{inc} is a set $I_s^\# \subseteq_s \mathbf{R}_{\text{inc}}$ such that $a \in_s I_s^\#$ and

$b \in_s I_s^\#$ imply $a + b \in_s I_s^\#$; $a \in_s I_s^\#$ and $r \in_s \mathbf{R}_{\text{inc}}$ imply $a \times r \in_s I_s^\#$, and $0_s \in I_s^\#$ and $1_s \notin I_s^\#$.

We define: (a) an strongly consistent equivalence relation $=_{I_s^\#, s}$ on \mathbf{R}_{inc} by $r =_{I_s^\#} s$ if and only if $r - s \in_s I_s^\#$.

(b) an weakly inconsistent equivalence relation $=_{I_s^\#, w}$ on \mathbf{R}_{inc} by $r =_{I_s^\#, w} s$ if and only if $r - s \in_w I_s^\#$.

We let $\mathbf{R}_{\text{inc}}/I_s^\#$ be the set of $r \in_s \mathbf{R}_{\text{inc}}$ such that r is the $<_{s,\mathbb{N}^\#}$ -minimum element of its equivalence class under $=_{I_s^\#, s}$. Thus $\mathbf{R}_{\text{inc}}/I_s^\#$ consists of one

element of each $=_{I_s^\#, s}$ -equivalence class of elements of \mathbf{R}_{inc} . With the

appropriate operations, $\mathbf{R}/I_s^\#$ becomes a countable commutative ring, the

quotient ring of \mathbf{R}_{inc} by $I_s^\#$. The ideal $I_s^\#$ is said to be prime if $\mathbf{R}_{\text{inc}}/I_s^\#$ is an

integral domain, and maximal if $\mathbf{R}_{\text{inc}}/I_s^\#$ is a field.

Next we indicate how some basic concepts and results of analysis and topology can be developed within $\mathbf{Z}_2^\#$.

Definition 3.2.17. Within $\mathbf{Z}_2^\#$, a s -paracomplete separable paraconsistent

metric space is a nonempty set $A \subseteq_s \mathbb{N}^\#$ together with a function

$d_s : A \times_s A \rightarrow_s \mathbb{R}^\#$ satisfying $a =_s a \rightarrow d_s(a, a) =_s 0_s$,

$0_s \leq_s d_s(a, b) =_s d_s(b, a)$,

and $d_s(a, c) \leq_s d_s(a, b) + d_s(b, c)$ for all $a, b, c \in_s A$.

(Formally, d_s is a strictly consistent sequence of paraconsistent real numbers, indexed by $A \times_s A$.)

We define a point of the s -paracomplete separable paraconsistent metric

space \widehat{A} to be a sequence $x =_s \langle a_n : n \in_s \mathbb{N}^\# \rangle_s$, $a_n \in_s A$, satisfying

$\forall \varepsilon (\varepsilon \in_s \mathbb{R}^\#) (0_s <_s \varepsilon$

The pseudometric d_s is extended from A to $(\widehat{A})_s$ by

$$d_s(x, y) = s\text{-} \lim_{n \rightarrow \infty} d_s(a_n, b_n) \quad (3.2.32)$$

where $x =_s \langle a_n : n \in \mathbb{N}^\# \rangle_s$ and $y =_s \langle b_n : n \in_s \mathbb{N}^\# \rangle_s$. We write $x =_s y$

if and only if $d_s(x, y) =_s 0_s$. For example, $\mathbb{R}^\# =_s \widehat{\mathbb{Q}^\#}$ under the metric $d_s(q, q') =_s |q - q'|_s$.

Definition 3.2.18. Within $\mathbf{Z}_2^\#$, a *weakly paracomplete* (w -paracomplete) separable paraconsistent metric space is a nonempty set $A \subseteq_s \mathbb{N}^\#$

together with a function $d_w : A \times_s A \rightarrow_s \mathbb{R}^\#$ satisfying
 $a =_s a \rightarrow d_w(a, a) =_s 0_s, a =_w a \rightarrow d_w(a, a) =_w 0_w,$
 $0_w \leq_w d_w(a, b) =_w d_w(b, a), \text{ and } d_w(a, c) \leq_w d_w(a, b) + d_w(b, c) \text{ for all } a, b, c \in_s A.$

(Formally, d_w is a consistent sequence of paraconsistent real numbers, indexed by $A \times_s A$.)

We define a point of the w -paracomplete separable paraconsistent metric

space $(\widehat{A})_w$ to be a sequence $x =_s \langle a_n : n \in_s \mathbb{N}^\# \rangle_s, a_n \in_s A$, satisfying

$$\forall \varepsilon (\varepsilon \in_s \mathbb{R}^\#) (0_w <_w \varepsilon \rightarrow \exists m \forall n (m <_s n \rightarrow d_w(a_m, a_n) <_w \varepsilon)). \quad (3.2.33)$$

The pseudometric d_w is extended from A to $(\widehat{A})_w$ by

$$d_w(x, y) = w\text{-} \lim_{n \rightarrow \infty} d_w(a_n, b_n) \quad (3.2.34)$$

where $x =_s \langle a_n : n \in \mathbb{N}^\# \rangle_s$ and $y =_s \langle b_n : n \in_s \mathbb{N}^\# \rangle_s$. We write

- (a) $x =_s y$ if and only if $d_s(x, y) =_s 0_s$,
- (b) $x =_w y$ if and only if $d_w(x, y) =_w 0_w$.

For example, $\mathbb{R}_w^\# =_s \widehat{(\mathbb{Q}^\#)}_w$ under the metric $d_w(q, q') =_s |q - q'|_w$.

Definition 3.2.19. Within $\mathbf{Z}_2^\#$, a *weakly paracomplete* with $\text{rang}= n, n \in \omega$

($\{w, (n)\}$ -paracomplete) separable paraconsistent metric space is a nonempty set $A \subseteq_s \mathbb{N}^\#$ together with a function $d_w : A \times_s A \rightarrow_s \mathbb{R}^\#$

satisfying:

$a =_s a \rightarrow d_{w,(n)}(a, a) =_s 0_s, a =_{w,(n)} a \rightarrow d_{w,(n)}(a, a) =_{w,(n)} 0_{w,(n)},$
 $0_{w,(n)} \leq_{w,(n)} d_{w,(n)}(a, b) =_{w,(n)} d_{w,(n)}(b, a), \text{ and}$
 $d_{w,(n)}(a, c) \leq_{w,(n)} d_{w,(n)}(a, b) + d_{w,(n)}(b, c)$
for all $a, b, c \in_s A$.

(Formally, $d_{w,(n)}$ is a consistent sequence of paraconsistent real numbers, indexed by $A \times_s A$.)

We define a point of the $\{w, (n)\}$ -paracomplete separable paraconsistent

metric space $(\widehat{A})_{w,(n)}$ to be a sequence $x =_s \langle a_n : n \in_s \mathbb{N}^\# \rangle_s, a_n \in_s A$,

satisfying:

$$\forall \varepsilon (\varepsilon \in_s \mathbb{R}^\#) (0_{w,(n)} <_{w,(n)} \varepsilon \rightarrow \exists m \forall n (m <_s n \rightarrow d_{w,(n)}(a_m, a_n) <_{w,(n)} \varepsilon)). \quad (3.2.35)$$

Definition 3.2.20. Within $\mathbf{Z}_2^\#$, a *strictly paracomplete* with $\text{rang}= n, n \in \omega$

($\{w, [n]\}$ -paracomplete) separable paraconsistent metric space is a nonempty set $A \subseteq_s \mathbb{N}^\#$ together with a function $d_w : A \times_s A \rightarrow_s \mathbb{R}^\#$ satisfying:

$a =_s a \rightarrow d_{w,[n]}(a, a) =_s 0_s, a =_{w,[n]} a \rightarrow d_{w,[n]}(a, a) =_{w,[n]} 0_{w,[n]},$

$0_{w,[n]} \leq_{w,[n]} d_{w,[n]}(a, b) =_{w,[n]} d_{w,[n]}(b, a)$, and
 $d_{w,[n]}(a, c) \leq_{w,[n]} d_{w,[n]}(a, b) + d_{w,[n]}(b, c)$
 for all $a, b, c \in_s A$.

(Formally, $d_{w,[n]}$ is a consistent sequence of paraconsistent real numbers, indexed by $A \times_s A$.)

We define a point of the $\{w, [n]\}$ -paracomplete separable paraconsistent

metric space $(\widehat{A})_{w,[n]}$ to be a sequence $x =_s \langle a_n : n \in_s \mathbb{N}^\# \rangle_s$, $a_n \in_s A$,

satisfying:

$$\forall \varepsilon (\varepsilon \in_s \mathbb{R}^\#) (0_{w,[n]} <_{w,[n]} \varepsilon \rightarrow \exists m \forall n (m <_s n \rightarrow d_{w,[n]}(a_m, a_n) <_{w,[n]} \varepsilon)). \quad (3.2.3)$$

Definition 3.2.21. (paraconsistent s -continuous functions). Within $\mathbf{Z}_2^\#$, if

\widehat{A}

and \widehat{B} are complete separable paraconsistent metric spaces, a

paraconsistent s -continuous function $\phi : \widehat{A} \rightarrow_s \widehat{B}$ is a set

$\Phi_s \subseteq_s A \times_s (\mathbb{Q}^\#)^+ \times_s B \times_s (\mathbb{Q}^\#)^+$ satisfying the following coherence

$$1. [(a, r, b, s)_s \in_s \Phi_s] \wedge [(a, r, b', s')_s \in_s \Phi_s] \dashrightarrow d_s(b, b') <_s s + s';$$

conditions:

$$2. [(a, r, b, s)_s \in_s \Phi_s] \wedge [d_s(b, b') + s <_s s'] \dashrightarrow (a, r, b', s')_s \in_s \Phi_s \quad (3.2.36)$$

$$3. [(a, r, b, s)_s \in_s \Phi_s] \wedge [d_s(a, a') + r' <_s r] \dashrightarrow (a', r', b, s)_s \in_s \Phi_s$$

Definition 3.2.22. (paraconsistent w -paracontinuous functions). Within

$\mathbf{Z}_2^\#$,

if \widehat{A} and \widehat{B} are w -paracomplete separable paraconsistent metric spaces, a

paraconsistent w -paracontinuous function $\phi : \widehat{A} \rightarrow_s \widehat{B}$ is a set

$\Phi_w \subseteq_w A \times_s (\mathbb{Q}^\#)^+ \times_s B \times_s (\mathbb{Q}^\#)^+$ satisfying the following

$$1. [(a, r, b, s)_w \in_w \Phi_w] \wedge [(a, r, b', s')_w \in_w \Phi_w] \dashrightarrow d_w(b, b') <_w s + s';$$

coherence conditions:

$$2. [(a, r, b, s)_w \in_w \Phi_w] \wedge [d_w(b, b') + s <_w s'] \dashrightarrow (a, r, b', s')_w \in_w \Phi_w$$

$$3. [(a, r, b, s)_w \in_w \Phi_w] \wedge [d_w(a, a') + r' <_w r] \dashrightarrow (a', r', b, s)_w \in_w \Phi_w$$

Definition 3.2.23. (paraconsistent $\{w, (n)\}$ -paracontinuous functions).

Within $\mathbf{Z}_2^\#$, if \widehat{A} and \widehat{B} are $\{w, (n)\}$ -paracomplete separable paraconsistent

metric spaces, a paraconsistent $\{w, (n)\}$ -paracontinuous function $\phi : \widehat{A} \rightarrow_s$

\widehat{B}

is a set $\Phi_w \subseteq_w A \times_s (\mathbb{Q}^\#)^+ \times_s B \times_s (\mathbb{Q}^\#)^+$ satisfying the following coherence

1. $[(a, r, b, s)_{w,(n)} \in_{w,(n)} \Phi_{w,(n)}] \wedge [(a, r, b', s')_{w,(n)} \in_{w,(n)} \Phi_{w,(n)}] \dashrightarrow$
 $\rightarrow d_{w,(n)}(b, b') <_{w,(n)} s + s';$
2. $[(a, r, b, s)_{w,(n)} \in_{w,(n)} \Phi_{w,(n)}] \wedge [d_{w,(n)}(b, b') + s <_{w,(n)} s'] \dashrightarrow$
 $\rightarrow (a, r, b', s')_{w,(n)} \in_{w,(n)} \Phi_{w,(n)}$
3. $[(a, r, b, s)_{w,(n)} \in_{w,(n)} \Phi_{w,(n)}] \wedge [d_{w,(n)}(a, a') + r' <_{w,(n)} r] \dashrightarrow$
 $\rightarrow (a', r', b, s)_{w,(n)} \in_{w,(n)} \Phi_{w,(n)}$

Definition 3.2.24. (paraconsistent $\{w, [n]\}$ -paracontinuous functions).

Within $\mathbf{Z}_2^\#$, if \widehat{A} and \widehat{B} are $\{w, [n]\}$ -paracomplete separable paraconsistent metric spaces, a paraconsistent $\{w, [n]\}$ -paracontinuous function $\phi : \widehat{A} \rightarrow_s \widehat{B}$ is a set $\Phi_w \subseteq_w A \times_s (\mathbb{Q}^\#)^+ \times_s B \times_s (\mathbb{Q}^\#)^+$ satisfying the following coherence

1. $[(a, r, b, s)_{w,[n]} \in_{w,[n]} \Phi_{w,[n]}] \wedge [(a, r, b', s')_{w,[n]} \in_{w,[n]} \Phi_{w,[n]}] \dashrightarrow$
 $\rightarrow d_{w,[n]}(b, b') <_{w,[n]} s + s';$
2. $[(a, r, b, s)_{w,[n]} \in_{w,[n]} \Phi_{w,[n]}] \wedge [d_{w,[n]}(b, b') + s <_{w,[n]} s'] \dashrightarrow$
 $\rightarrow (a, r, b', s')_{w,[n]} \in_{w,[n]} \Phi_{w,[n]}$
3. $[(a, r, b, s)_{w,[n]} \in_{w,[n]} \Phi_{w,[n]}] \wedge [d_{w,[n]}(a, a') + r' <_{w,[n]} r] \dashrightarrow$
 $\rightarrow (a', r', b, s)_{w,[n]} \in_{w,[n]} \Phi_{w,[n]}$

Definition 3.2.25. (paraconsistent s -open sets). Within $\mathbf{Z}_2^\#$, let \widehat{A} be a s -paracomplete separable paraconsistent metric space. A (code for an) strictly open set (s -open set) in \widehat{A} is any set $U \subseteq_s \widehat{A} \times (\mathbb{Q}^\#)^+$. For $x \in_s \widehat{A}$ we

write $x \in_s U$ if and only if $d_s(x, a) <_s r$ for some $(a, r)_s \in_s U$.

Definition 3.2.26. (paraconsistent w -open sets). Within $\mathbf{Z}_2^\#$, let \widehat{A} be a w -paracomplete separable paraconsistent metric space. A (code for an) weakly open set (w -open set) in \widehat{A} is any set $U \subseteq_w \widehat{A} \times (\mathbb{Q}^\#)^+$. For $x \in_w \widehat{A}$ we write $x \in_w U$ if and only if $d_w(x, a) <_w r$ for some $(a, r)_w \in_w U$.

Definition 3.2.27. (paraconsistent $\{w, (n)\}$ -open sets). Within $\mathbf{Z}_2^\#$, let \widehat{A} be a

$\{w, (n)\}$ -paracomplete separable paraconsistent metric space.

A (code for an) weakly open with $\text{rang} = n, n \in \omega$ set ($\{w, (n)\}$ -open set) in \widehat{A}

is any set $U \subseteq_{w, (n)} \widehat{A} \times (\mathbb{Q}^\#)^+$. For $x \in_{w, (n)} \widehat{A}$ we write $x \in_{w, (n)} U$ if and only

if $d_{w, (n)}(x, a) <_{w, (n)} r$ for some $(a, r)_{w, (n)} \in_{w, (n)} U$.

Definition 3.2.28. (paraconsistent $\{w, [n]\}$ -open sets). Within $\mathbf{Z}_2^\#$, let \widehat{A} be a

$\{w, [n]\}$ -paracomplete separable paraconsistent metric space.

A (code for an) weakly open with $\text{rang} = n, n \in \omega$ set ($\{w, [n]\}$ -open set) in \widehat{A}

is any set $U \subseteq_{w, [n]} \widehat{A} \times (\mathbb{Q}^\#)^+$. For $x \in_{w, [n]} \widehat{A}$ we write $x \in_{w, [n]} U$ if and only

if $d_{w, [n]}(x, a) <_{w, [n]} r$ for some $(a, r)_{w, [n]} \in_{w, [n]} U$.

Definition 3.2.29. A separable paraconsistent Banach \mathbf{s} -space is defined

within $\mathbf{Z}_2^\#$ to be a paracomplete separable metric space \widehat{A} arising from a

countable \mathbf{s} -pseudonormed vector space \widehat{A} over the paraconsistent field $\mathbb{Q}^\#$.

Example 3.2.1. (a) With the \mathbf{s} -metric

$$d_{\mathbf{s}}(f, g) =_{\mathbf{s}} \sup_{0_{\mathbf{s}} \leq_{\mathbf{s}} x \leq_{\mathbf{s}} 1_{\mathbf{s}}} |f_1(x) - g_1(x)|$$

we have \mathbf{s} -paracomplete separable paraconsistent metric

space $\widehat{A} \triangleq \widehat{C_{\mathbf{s}}[0_{\mathbf{s}}, 1_{\mathbf{s}}]_{\mathbf{s}}}$, where $C_{\mathbf{s}}[0_{\mathbf{s}}, 1_{\mathbf{s}}]_{\mathbf{s}}$ is a paraconsistent linear space paraconsistent \mathbf{s} -continuous functions $f_{\mathbf{s}} : [0_{\mathbf{s}}, 1_{\mathbf{s}}]_{\mathbf{s}} \rightarrow_{\mathbf{s}} \mathbb{R}^\#$.

(b) With the w -metric

$$d_w(f, g) =_w \sup_{0_w \leq_w x \leq_w 1_w} |f_1(x) - g_1(x)|_w \quad (3.2.41)$$

we have w -paracomplete separable paraconsistent metric

space $\widehat{A} \triangleq \widehat{C_w[0_w, 1_w]}_w$, where $C_w[0_w, 1_w]_w$ is a paraconsistent linear space paraconsistent w -continuous functions $f_w : [0_w, 1_w]_w \rightarrow_w \mathbb{R}^\#$.

9 IV.Berry's and Richard's inconsistent numbers within $\mathbf{Z}_2^\#$.

10 IV.1.Hierarchy Berry's inconsistent numbers $\mathbf{B}_n^{w,(m)}$.

Suppose that $F_{n_1}^w(n, X) \in L_2^\#$ is a well-formed formula of second-order arithmetic $\mathbf{Z}_2^\#$, i.e. formula which is arithmetical, which has one free set variable X and one free individual variable n . Suppose that $g(\exists X F_{n_1}^w(x, X)) \leq \mathbf{k}$, where $g(\exists X F_{n_1}^w(x, X))$ is a corresponding Gödel number. Let be $A_{\mathbf{k}}^w, \mathbf{k} \in \mathbb{N}$ the set of all positive weakly inconsistent integers $\bar{n} \in_s \mathbb{N}_w^\#$ which can be defined within $\mathbf{Z}_2^\#$ (in weak inconsistent sense) under corresponding well-formed formula $F_{\bar{n}_1(\bar{n})}(x, X)$, i.e. $\exists X_{\bar{n}} \forall m [F_{\bar{n}_1(\bar{m})}^w(\bar{m}, X_{\bar{n}}) \rightarrow \bar{m} =_w \bar{n}]$, hence $\bar{n} \in A_{\mathbf{k}}^w \longleftrightarrow \exists X_{\bar{n}} F_{\bar{n}_1(\bar{n})}^w(\bar{n}, X_{\bar{n}})$.

$$\text{Thus } \begin{aligned} \forall n_{n \in_s \mathbb{N}_w^\#} [n \in A_{\mathbf{k}}^w \longleftrightarrow \exists X_n F_{n_1(n)}(n, X_n)], \\ \text{where } g(\exists X_n F_{n_1}(x, X_n)) \leq \mathbf{k}, \end{aligned} \quad (4.1.1)$$

Since there are only finitely many of these \bar{n} , there must be a smallest (relative to $<_w$) positive integer $\mathbf{B}_{\mathbf{k}}^w \in_w \mathbb{N}_w^\# \setminus_w A_{\mathbf{k}}^w$ that does not belong to

$A_{\mathbf{k}}^w$. But we just defined $\mathbf{B}_{\mathbf{k}}^w$ in under corresponding well-formed formula

$$\breve{F}_{\bar{n}_1(\mathbf{B}_{\mathbf{k}})}(\mathbf{B}_{\mathbf{k}}^w, \mathbf{k})$$

Hence for a sufficiently Large \mathbf{k} such that: $g(\breve{F}_{\bar{n}_1(\mathbf{B}_{\mathbf{k}})}(\mathbf{B}_{\mathbf{k}}^w, A_{\mathbf{k}})) \leq \mathbf{k}$ we obtain:

$$(\mathbf{B}_{\mathbf{k}}^w \in_w A_{\mathbf{k}}^w) \wedge (\mathbf{B}_{\mathbf{k}}^w \notin_w A_{\mathbf{k}}^w). \quad (4.1.3)$$

Theorem.4.1.1. Paraconsistent set $A_{\mathbf{k}}^w \subset_w \mathbb{N}_w^\#$ which was defined above it a strictly \in -inconsistent set with **rank** ≥ 0 .

11 IV.2.Hierarchy Richard's inconsistent numbers $\mathfrak{R}_n^{w,(m)}$.

Let be $q_n \in_s \mathbb{Q}_w^\#$ paraconsistent rational number with corresponding decimal representation $q_n =_s \{0, q_n(1_w) q_n(2_w) \dots q_n(i) \dots q_n(n)\}, n \in_s \mathbb{N}_w^\#$,
 $q_n(i) =_w 0_w \vee 1_w \vee 2_w \vee 3_w \vee 4_w \vee 5_w \vee 6_w \vee 7_w \vee 8_w \vee 9_w$,

$i \leq_w n, x_k^w =_s \langle q_n^k \rangle =_s \langle q_n^k : n \in \mathbb{N}, q_n^k =_s \{0, q_n^k(1) q_n^k(2) \dots q_n^k(n)\} \rangle_s \in_s \mathbb{R}_w^\#$, $k \in \mathbb{N}$ is a paraconsistent real number which can be defined in weak inconsistent sense under corresponding well-formed formula (of second-order arithmetic $\mathbf{Z}_2^\#$) $F_k(x, X)$, i.e. $\forall q (q \in_s \mathbb{Q}_w^\#) [q \in_w \langle q_n^k \rangle \leftrightarrow \exists X F_k(q, X)]$.

Definition. 4.2.1. We denote paraconsistent real number x_k^w as k -th Richard's weakly inconsistent real number.

Let us consider Richard's real number $\mathfrak{R}_p^w =_s \langle \mathfrak{R}_n^p : n \in \mathbb{N} \rangle$ such that

Suppose that $q_p^p(p) \neq_w 1$, hence $\mathfrak{R}_p^p(p) =_w 1_w$. Thus $\mathfrak{R}_p^p(p) \neq_w q_p^p(p) \rightarrow \mathfrak{R}_p \neq_w x_p$. Suppose that $q_p^p(p) =_w 1_w$, hence $\mathfrak{R}_p^p(p) =_w 0_w$. Thus $\mathfrak{R}_p^p(p) \neq_w q_p^p(p) \rightarrow \mathfrak{R}_p \neq_w x_p$.

Hence for any Richard's real number x_k one obtain the contradiction

Theorem.4.2.1. Paraconsistent set $\{x_k^w : k \in \mathbb{N}\}_s \subset_w \mathbb{R}_w^\#$ containing the all Richard's weakly inconsistent real numbers which was defined above it a strictly \in -inconsistent set with **rank** ≥ 0 .

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